



**Putting Order
into
Classifications and Universality**
AUTOMATA 2010

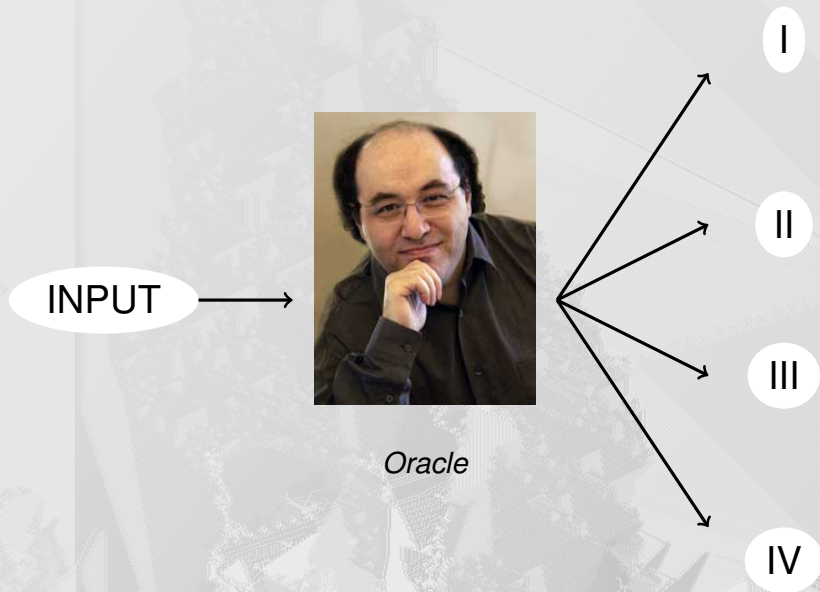
Guillaume Theyssier

LAMA lab. (CNRS, Université de Savoie, France)

June 14th, 2010

- Banks (1970)
- Albert and Čulík (1987)
- Martin I (1993)
- Durand and Róka (1996)
- Durand-Lose (1997)
- Mazoyer and Rapaport (1998)
- Ollinger (2002)
- T (2005)

What is a classification?



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► **Classical approach:**

- 1** define *a priori* a list of properties or **classes**
- 2** being **similar** \Leftrightarrow satisfying the same properties

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► This talk:

- 1 define a relation of **similarity** or equivalence between CA
- 2 equivalence **classes** are obtained *a posteriori*

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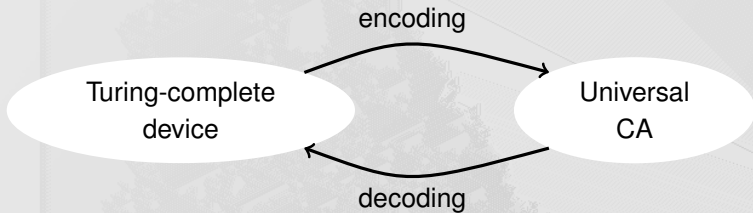
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Key points:

- no external tools used to define properties
- infinitely many classes

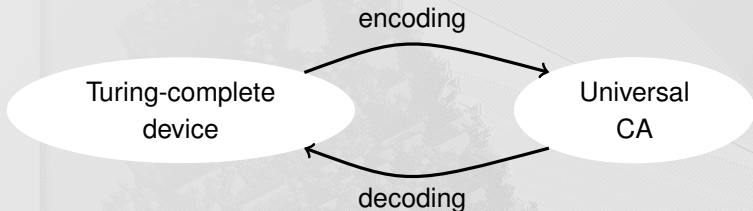
What is universality?

► Classical approach:



What is universality?

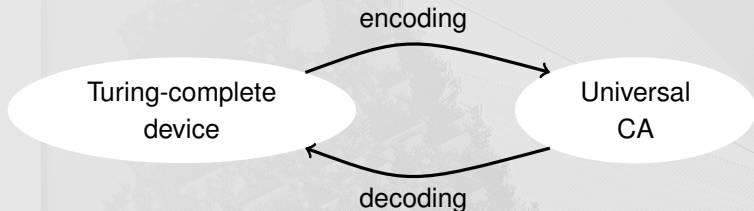
► Classical approach:



- encoding/decoding/halt problem
- only a positive definition

What is universality?

► Classical approach:



- encoding/decoding/halt problem
- only a positive definition

► This talk:

- 1 define a notion of **simulation** between CA
- 2 **intrinsic universality** $\stackrel{def}{=} ability to simulate all CA$

Putting Order into Classifications and Universality

- ▶ *1 solution for 2 problems*



Putting Order into Classifications and Universality

► *1 solution for 2 problems*

1 define a **pre-order** on CA: \preceq

(reflexive and transitive relation)

Putting Order into Classifications and Universality

► 1 solution for 2 problems

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(reflexive and transitive relation)

2 classification:

- induced equivalence relation

$$F \sim G \stackrel{\text{def}}{\iff} F \preceq G \text{ and } G \preceq F$$

- topology of the pre-order

Putting Order into Classifications and Universality

► 1 solution for 2 problems

1 define a **pre-order** on CA: \preceq

(reflexive and transitive relation)

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- topology of the pre-order

3 universality:

$$F \text{ universal} \stackrel{\text{def}}{\iff} \forall G, G \preceq F$$

► *Ingredients:*

- 1 “local” comparison relations
 - subautomaton
 - factor
- 2 rescaling operations

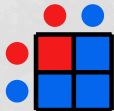
General idea

pre-order \equiv local comparison **up to** rescaling

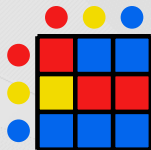
■ $F \sqsubseteq G$



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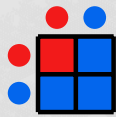


F

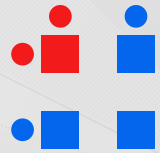


G

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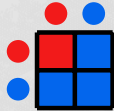
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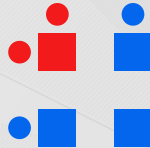
G

injection $\iota : Q_F \rightarrow Q_G$ with $\iota \circ F = G \circ \iota$

■ $F \sqsubseteq G$



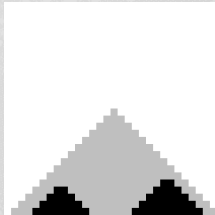
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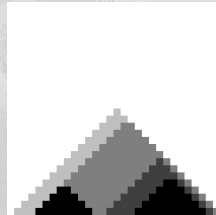
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■ example:



MAX with 3 states

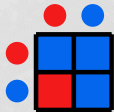
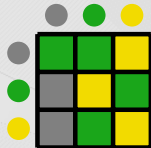


MAX with 5 states

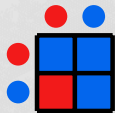
■ $F \trianglelefteq G$



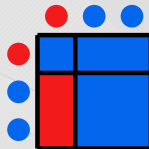
$$\blacksquare F \trianglelefteq G$$

 F  G

■ $F \trianglelefteq G$



F



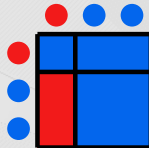
G

surjection $\pi : Q_G \rightarrow Q_F$ with $\pi \circ G = F \circ \pi$

■ $F \trianglelefteq G$



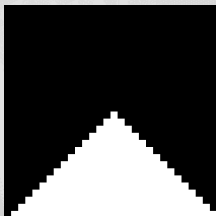
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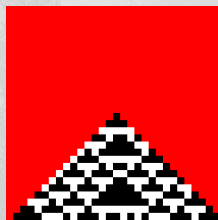
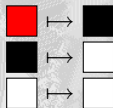
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■ example



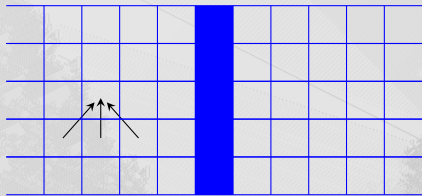
MAX with 2 states



54 + spreading state

Rescaling operations

► 3 parameters: $F \mapsto F\langle m, t, z \rangle$



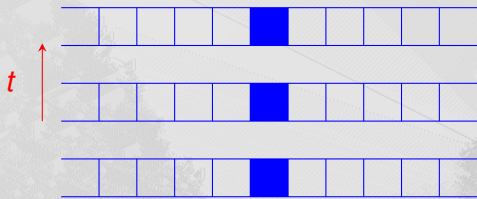
Global map

$$F\langle 1, 1, 0 \rangle = F$$

Rescaling operations

► 3 parameters: $F \mapsto F\langle m, t, z \rangle$

■ time



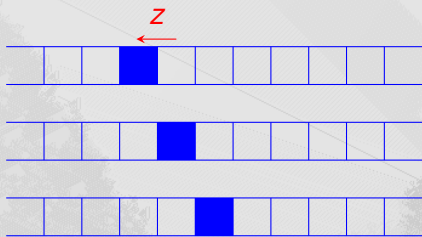
Global map

$$F\langle 1, t, 0 \rangle = F^t$$

Rescaling operations

► 3 parameters: $F \mapsto F\langle m, t, z \rangle$

- time
- shift



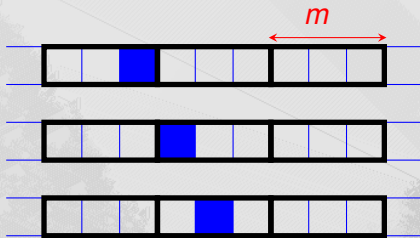
Global map

$$F\langle 1, t, z \rangle = \sigma_z \circ F^t$$

Rescaling operations

► 3 parameters: $F \mapsto F\langle m, t, z \rangle$

- time
- shift
- cell grouping



Global map

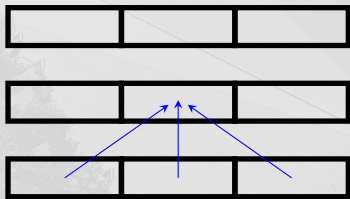
$$F\langle m, t, z \rangle = \mathbf{o}_m^{-1} \circ \sigma_z \circ F^t \circ \mathbf{o}_m$$

$$\mathbf{o}_m^{-1} : Q^{\mathbb{Z}} \rightarrow (Q^m)^{\mathbb{Z}} \text{ canonical bijection}$$

Rescaling operations

► 3 parameters: $F \mapsto F\langle m,t,z \rangle$

- time
- shift
- cell grouping

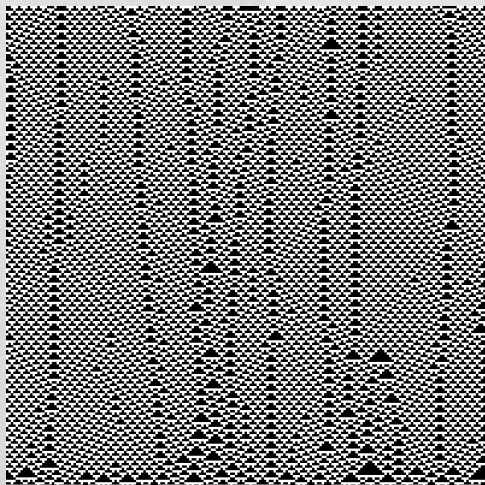


Fact

$F\langle m,t,z \rangle$ is a cellular automaton

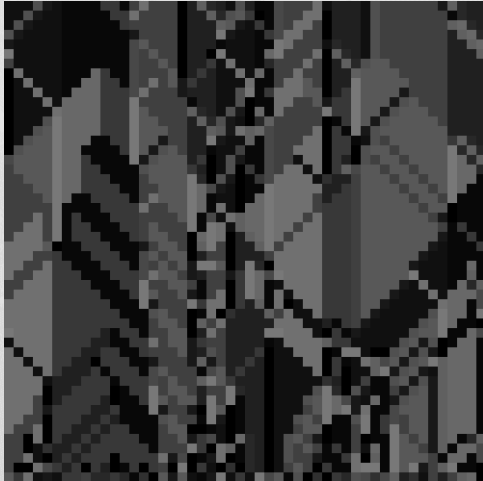
- with a possibly different alphabet
- with a possibly different radius

54

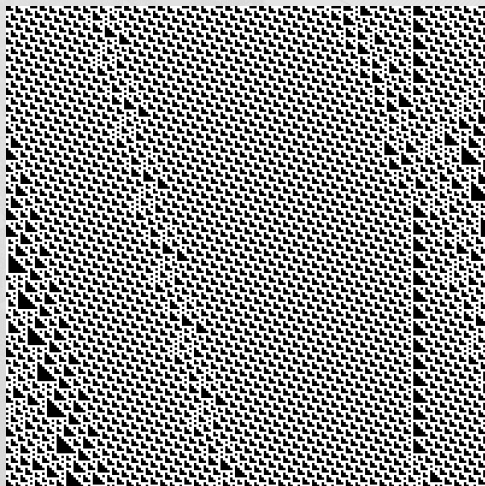


Rescaling examples

$54^{(4,4,0)}$



110



Rescaling examples

$110\langle 14,7,0\rangle$



3 pre-orders



■ **injective** simulation

$$F \preceq_i G \stackrel{\text{def}}{\iff} \exists \vec{p}_1, \vec{p}_2 : F \langle \vec{p}_1 \rangle \sqsubseteq G \langle \vec{p}_2 \rangle$$

■ injective simulation

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■ surjective simulation

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■ **surjective** simulation

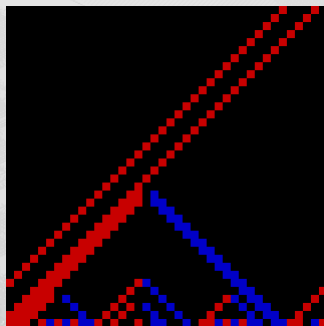
$$F \preceq_s G \stackrel{\text{def}}{\Leftrightarrow} \exists \vec{p}_1, \vec{p}_2 : F \langle \vec{p}_1 \rangle \trianglelefteq G \langle \vec{p}_2 \rangle$$

■ **mixed** simulation

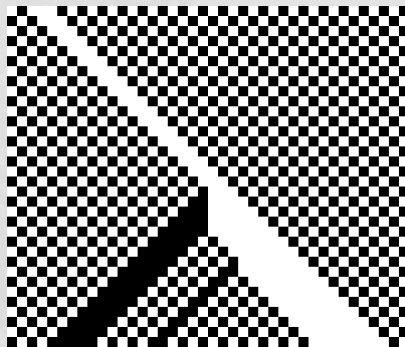
$$F \preceq_m G \stackrel{\text{def}}{\Leftrightarrow} \exists \vec{p}_1, \vec{p}_2 : F \langle \vec{p}_1 \rangle \trianglelefteq \sqsubseteq G \langle \vec{p}_2 \rangle$$



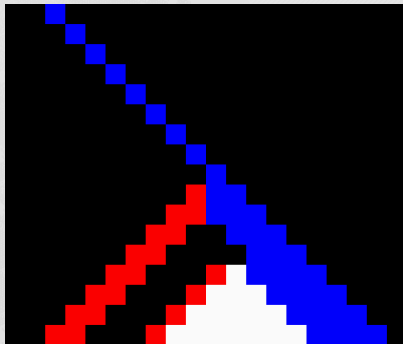
ECA 184



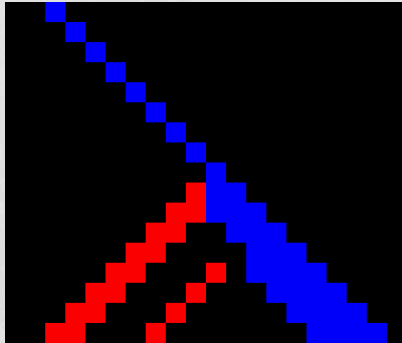
'Just Gliders'



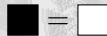
'Just Gliders' $\trianglelefteq 184^{(2,2)} \sim 184$

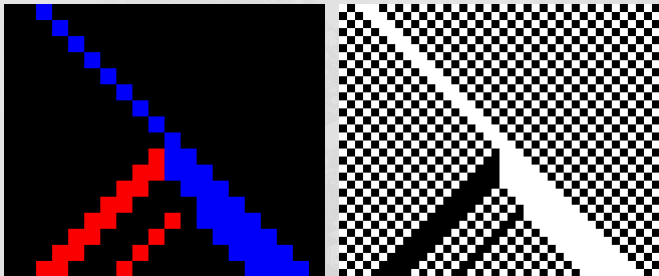


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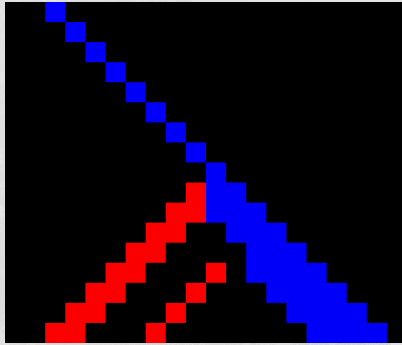


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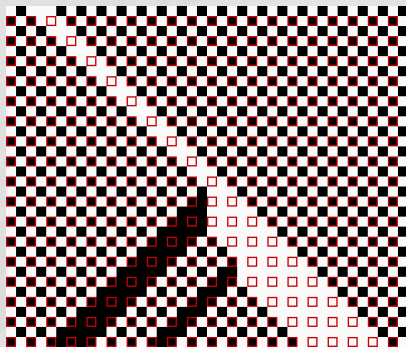




'Just Gliders' \preceq_s 184

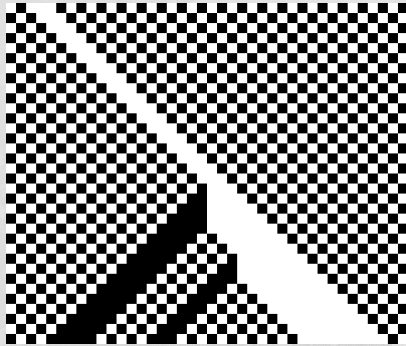


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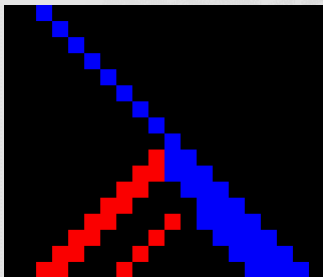


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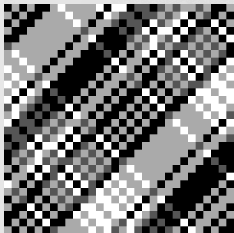
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'Just Gliders' \preceq_j 184

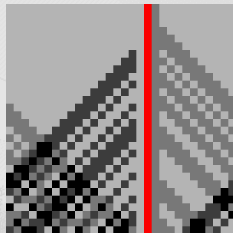
Separation





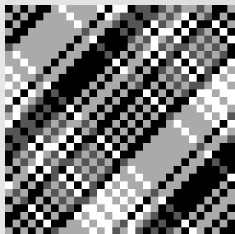
$$\sigma \times \sigma^{-1}$$

λ_i
 ~~λ_s~~
 λ_m



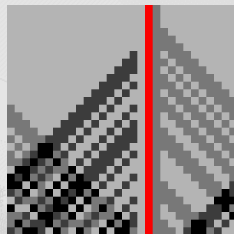
$$\sigma \times \sigma^{-1} + \text{wall state}$$

Separation

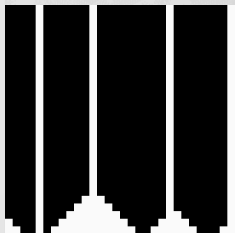


$\sigma \times \sigma^{-1}$

$\underbrace{\quad}_i$
 ~~$\underbrace{\quad}_s$~~
 $\underbrace{\quad}_m$

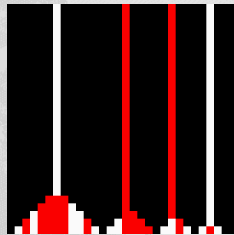


$\sigma \times \sigma^{-1} + \text{wall state}$



block reduction

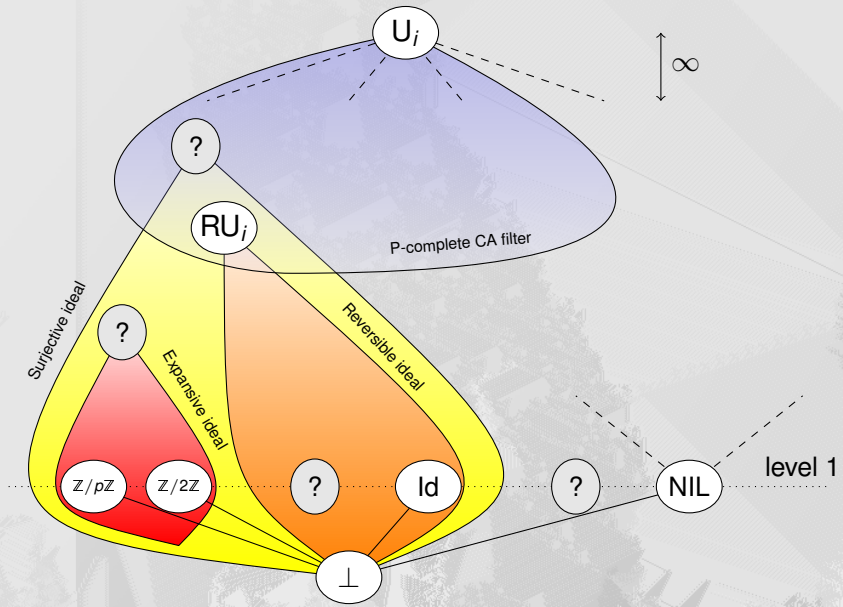
~~$\underbrace{\quad}_i$~~
 $\underbrace{\quad}_s$
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block reduction + parity test

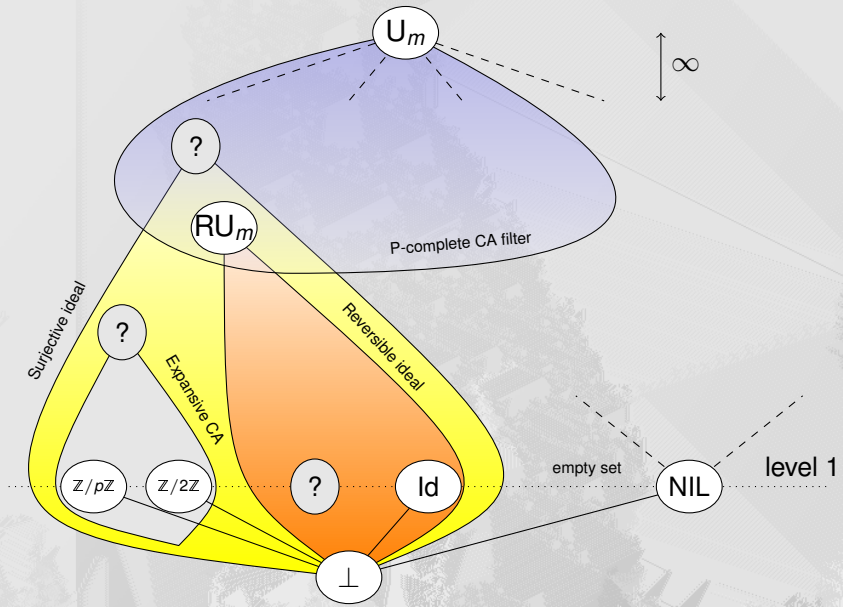
Pre-order 





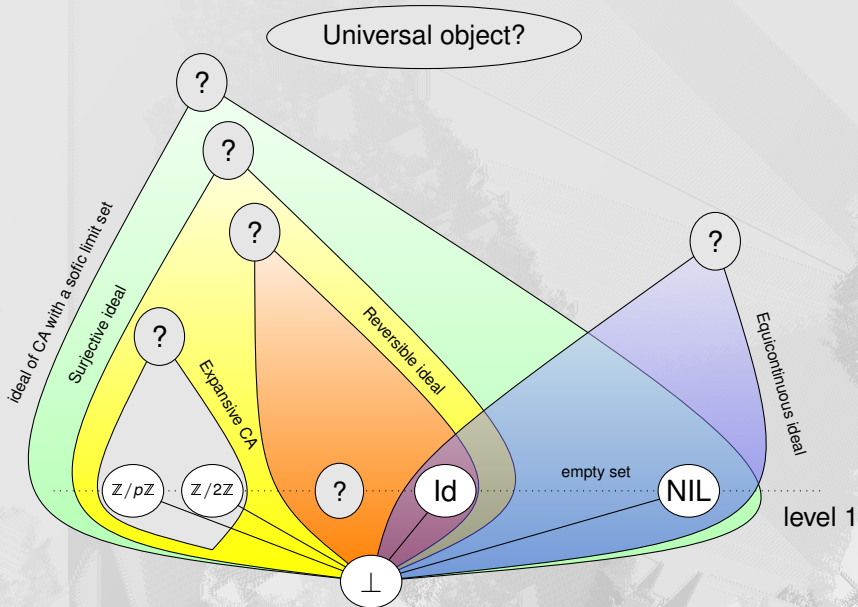
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intrinsic universality $\stackrel{\text{def}}{=} U_i \stackrel{\text{def}}{=} \{F : \forall G, G \preceq_i F\}$



N. Ollinger, “Universalities in Cellular Automata”,
Handbook of natural computing

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Theorem (Rapaport) — no real-time universality

If $F \in U_i$ then there is some G with $G^{\langle \vec{p} \rangle} \not\sqsubseteq F^{\langle t, t, z \rangle}, (\forall \vec{p}, t, z)$.

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Theorem (Ollinger) — strong universality

If $F \in U_i$ then for all G there is \vec{p} such that $G \sqsubseteq F^{\langle \vec{p} \rangle}$.

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Theorem (Ollinger) — strong universality

If $F \in U_i$ then for all G there is \vec{p} such that $G \sqsubseteq F^{\langle \vec{p} \rangle}$.

Exercise

If $F \in U_i$ and has a spreading state then it can simulate any G without using the spreading state.

Theorem (Ollinger)

Intrinsic universality is undecidable.

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Proposition

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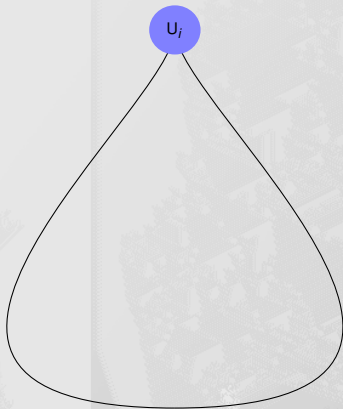
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Exercise

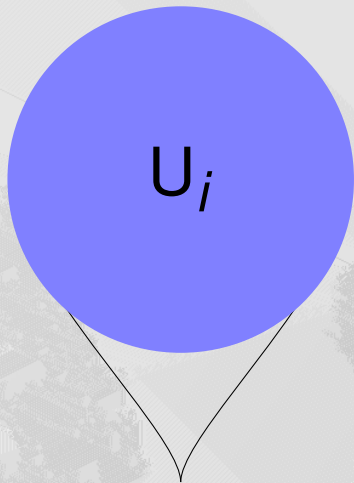
If F is not universal then there is a non-universal G with

$$F \preceq_i G \text{ but } G \not\preceq_i F$$

How Common is Intrinsic Universality?



or



How Common is Intrinsic Universality?

The quest for small intrinsically universal CA



How Common is Intrinsic Universality?

The quest for small intrinsically universal CA

► 2D

- Banks 1970: 2 states + von Neuman neighb.
- Conway 1970: Game of Life (2 states + Moore neighb.)

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The quest for small intrinsically universal CA

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- Banks 1970: 2 states + large radius
- Ollinger 2002: 6 states + radius 1
- Richard 2008: 4 states + radius 1 (U_i or U_m ?)

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Open problems

- $U_j \stackrel{?}{=} U_m$
- is 110 intrinsically universal?

How Common is Intrinsic Universality?

Adding syntactical constraints



How Common is Intrinsic Universality?

Adding syntactical constraints

► Definitions:

- F is **captive** if $\forall \vec{x}$:

$$f(x_1, \dots, x_k) \in \{x_1, \dots, x_k\}$$

- F is **multiset** if $\forall \vec{x}$ and \forall permutation π

$$f(\vec{x}) = f(\pi(\vec{x}))$$

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Theorem

There exists intrinsically universal CA in the following families:

- number conserving
- totalistic (\subseteq multiset)
- captive

How Common is Intrinsic Universality?

Density

- \mathcal{P} a **property** (a set of CA)
- \mathcal{F} a **family** (a set of CA)
- $\mathcal{F}_{n,r}$: CA from \mathcal{F} with n states and radius r

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$$d_n(\mathcal{P}/\mathcal{F}) = \lim_{n \rightarrow \infty} \frac{\#\mathcal{P} \cap \mathcal{F}_{n,r_0}}{\#\mathcal{F}_{n,r_0}}$$

$$d_r(\mathcal{P}/\mathcal{F}) = \lim_{r \rightarrow \infty} \frac{\#\mathcal{P} \cap \mathcal{F}_{n_0,r}}{\#\mathcal{F}_{n_0,r}}$$

How Common is Intrinsic Universality?

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$$d_n(\mathcal{P}/\mathcal{F}) = \lim_{n \rightarrow \infty} \frac{\#\mathcal{P} \cap \mathcal{F}_{n,r_0}}{\#\mathcal{F}_{n,r_0}}$$

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Theorem (Boyer,T.)

- $d_n(U_i/\text{captive CA}) = 1$
- $d_r(U_i/\text{multiset CA}) = 1$
- combination of both + other families (e.g. outer-totalistic)

How Common is Intrinsic Universality?

Density

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Open problem

- $d_n(U_i/CA)$?
- $d_r(U_i/CA)$?

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syntactical restriction \Rightarrow restriction on global behaviours?

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Theorem

There are recursive fair encodings from CA into captive CA, and from CA into multiset CA.

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What are equivalence classes without any captive and/or multiset CA?

Other universalities



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Open problem

Is there a “surjective-universal” CA in dimension 1?

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■ **q persistent state** $\stackrel{\text{def}}{\iff} f(*, q, *) = q$

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Open problem

Is there a universal CA for \preceq_s ?

- 1** more general simulations?
 - sub-systems induced by stable sub-shifts (SFT? sofic?)
 - factor with context (sliding block codes)
- 2** dimension change
 - constructions à la Toffoli?
 - sub-actions à la Hochman?
- 3** asynchronous CA

Don't Know What to Do During the World Cup?



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► My list of small open questions

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Thank you!