

Stabilizing SFT by Cellular Automata

Séminaire CANA

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- **[FMT21]:** “*Self-stabilisation of cellular automata on tilings*”
by **N. Fatès**, **I. Marcovici** and **S. Taati**

- **unpublished joint work** with **M. Delacourt**, **P. Guillon**, **N. Ollinger**

- **earlier discussions** with **A. Gajardo**, **D. Maldonado** and **S. Taati**

Subshift of finite type (SFT)

- A : finite **alphabet**
- $A^{\mathbb{Z}^d}$: d -dimensional **configurations**
- **patterns** : $p : D \rightarrow A$ with $D \subseteq \mathbb{Z}^d$ **finite**
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- in this talk (wlog), SFT using only domino domains:
 - 1D: $D = \{0, 1\}$
 - 2D: $D = \{(0, 0), (1, 0)\}$ and $D = \{(0, 0), (0, 1)\}$

Cellular automata (CA)

- defined by a **local map** $\lambda : A^D \rightarrow A$
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- **asymptotic** configurations: $x \stackrel{\infty}{\cong} y$ if $\{z : x(z) \neq y(z)\}$ is finite

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$x \in \Sigma$

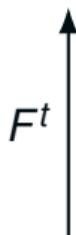
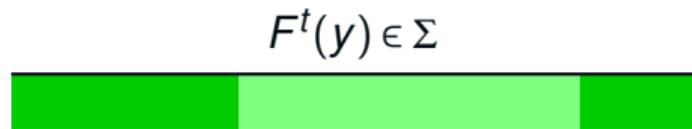
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The following 2D CA (Toom-majority rule) stabilizes Σ :

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Theorem (Gács-Kurdyumov-Levin-78)

The following 1D CA stabilizes Σ :

$$F(x)_z = \begin{cases} \text{maj}(x_z, x_{z+1}, x_{z+3}) & \text{if } x_z = 1 \\ \text{maj}(x_z, x_{z-1}, x_{z-3}) & \text{else.} \end{cases}$$



Local errors VS. distance to Σ

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- number of **local errors**

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$$d(y, \Sigma) = \min\{|D| : \exists x' \in \Sigma, \forall z \notin D, y(z) = x'(z)\}$$

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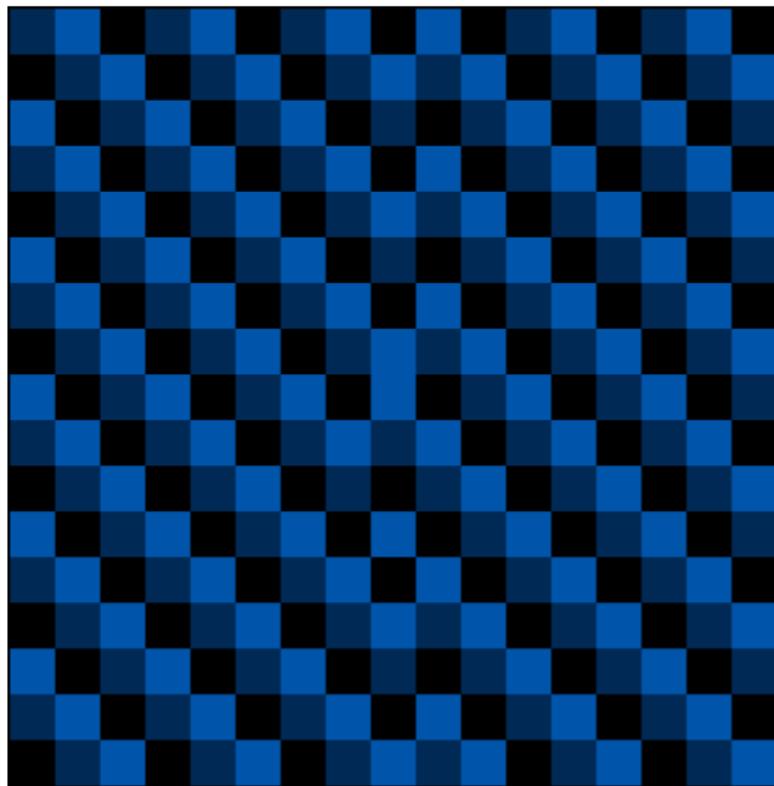
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-  we can have $\epsilon(y) \ll d(y, \Sigma)$
- a stabilizer must be able to **change arbitrary large patches of locally valid cells!**
- **stabilization time** as a function of $d(y, \Sigma)$

Example: $\Sigma = 3$ -colorings



Previous results from [FMT21]

Theorem [FMT21]

Any non-wandering 1D SFT is stabilizable in linear time.

- *intuition*: can be patched from left to right.

Theorem [FMT21]

Any NE-deterministic 2D SFT is stabilizable.

- *intuition*: can be reduced to a 1D patching process using the determinism.

New: a general result in 2D

Theorem 1

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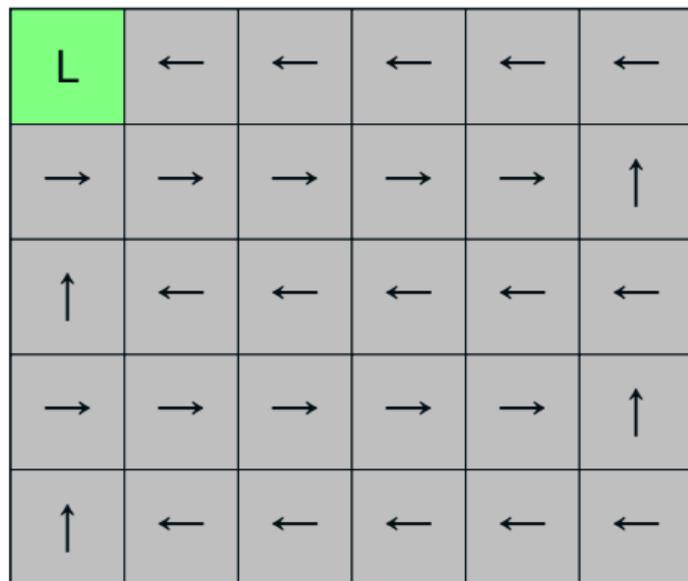
Main ingredients

- $B = A \cup A \times M$: cells in $A \times M$ are said **marked**
- **zone** = maximal 4-connected region of marked cells
- what the CA F do:
 - 1 **mark** any locally invalid cell
 - 2 **normalize** zones (to be defined in next slide)
 - 3 brute force **computation** inside normalized zone
 - if valid patch found → **remove the marks**
 - if not → **extend zone by 1 cell in each direction**

Normalized zones

| | | | | | |
|---|---|---|---|---|---|
| L | ← | ← | ← | ← | ← |
| → | → | → | → | → | ↑ |
| ↑ | ← | ← | ← | ← | ← |
| → | → | → | → | → | ↑ |
| ↑ | ← | ← | ← | ← | ← |

Normalized zones



- Turing computation inside: *try all possible A-filling with backtrack*
- no computation loop, even with bad initialization

- $\Sigma = 3$ -colorings in 2D
- a CA F with 69 states and radius 1 that stabilizes it

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- 4 impossible to have a constant number of zones ≥ 2 forever
- 5 when a single zone remains, it unmarks itself once large enough

δ -stabilizing under ϵ -perturbations

- $A = \{0, \dots, k-1\}$
- μ Bernoulli measure over $A^{\mathbb{Z}^d}$ with $\mu(0) \geq 1 - \epsilon$
- ρ a μ -**random configuration**
- ϵ -**perturbation** of x : $x + \rho = z \mapsto x(z) + \rho(z) \bmod k$

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F δ -stabilizes Σ under ϵ -perturbation if:

- $\forall x \in \Sigma, F(x) = x$
- $\forall x \in \Sigma$, for μ -almost all ρ :
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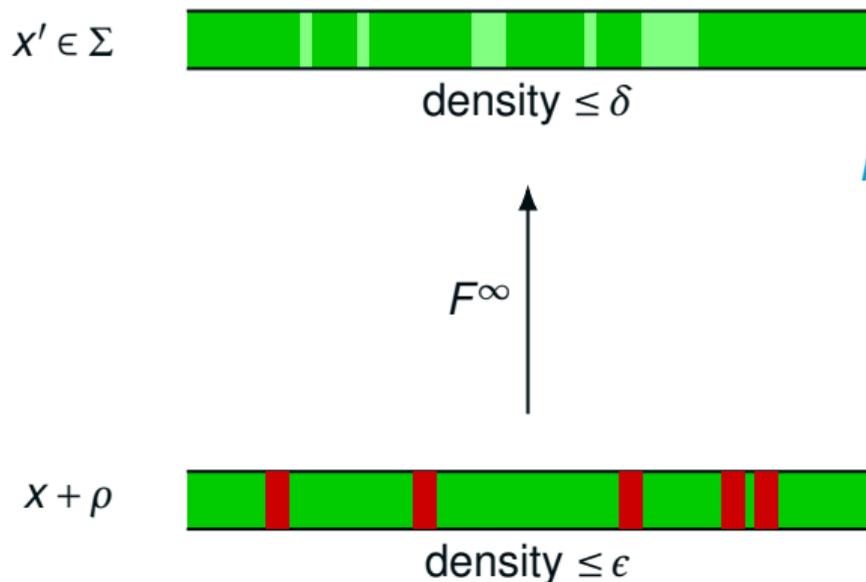
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If Σ can be stabilized in **linear time**, then

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Theorem [FMT21]

Unless $P=NP$, there is Σ in 2D which cannot be stabilized in polynomial time.

New: general linear time stabilization in 1D

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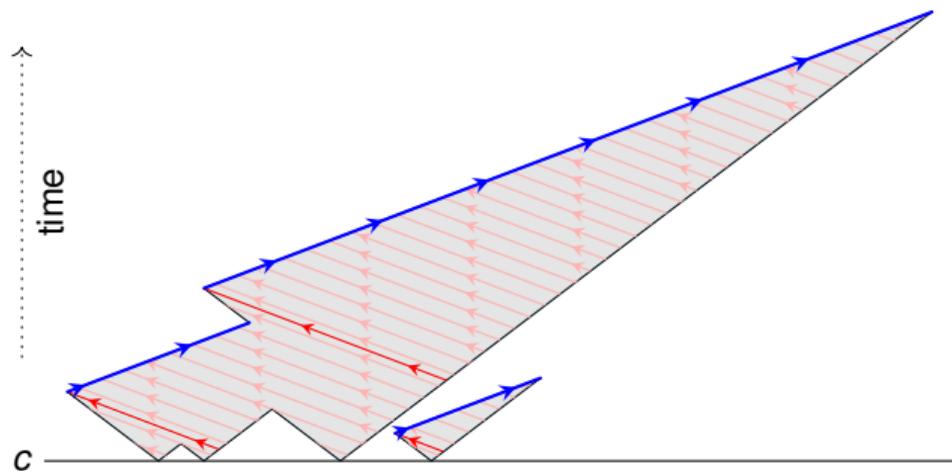
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Corollary

For any 1D SFT Σ , $\forall \delta, \exists \epsilon$ such that Σ can be δ -stabilized against ϵ -perturbations.

Construction idea

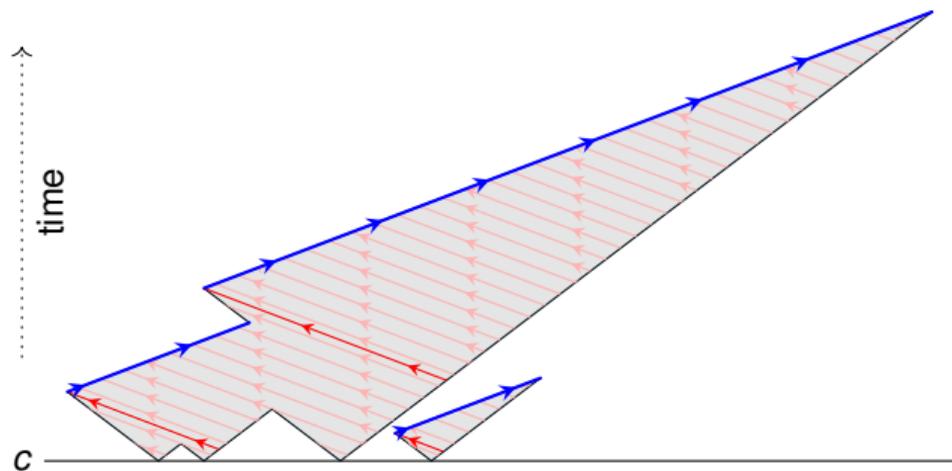
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-  exact definition & analysis a bit subtle...

Work in progress & future work

- if context-filling problem of Σ is **polytime** then stabilization in **polytime**
- remove extra symbols (*i.e.* $B = A$)
- which SFT have a polytime context-filling problem?
- what about random perturbation in 2D? e.g. with 3-colouring?