

Directional Dynamics of Cellular Automata

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LIF / LIF / LATP / LAMA
(CNRS, Université de Provence) $\times 3$, (CNRS, Université de Savoie)

September 1st 2010

Overview of the talk

- 1 Cellular Automata**
- 2 Topological Dynamics**
- 3 Directional Dynamics**

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1 Cellular Automata

2 Topological Dynamics

3 Directional Dynamics

Cellular automata

Definition

► Syntactical object:

- Q state set,
- $V = \{v_1, \dots, v_k\} \subseteq \mathbb{Z}$ neighborhood,
- $f : Q^V \rightarrow Q$ **local transition rule**

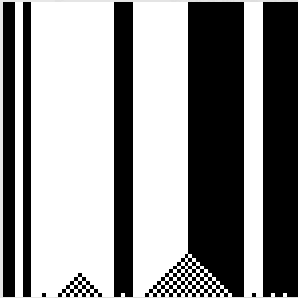
► Associated behavior:

- $Q^{\mathbb{Z}}$: set of configurations
- $F : Q^{\mathbb{Z}} \rightarrow Q^{\mathbb{Z}}$ **global transition rule** defined by

$$F(x)_z = f(x_{z+v_1}, \dots, x_{z+v_k})$$

Cellular automata

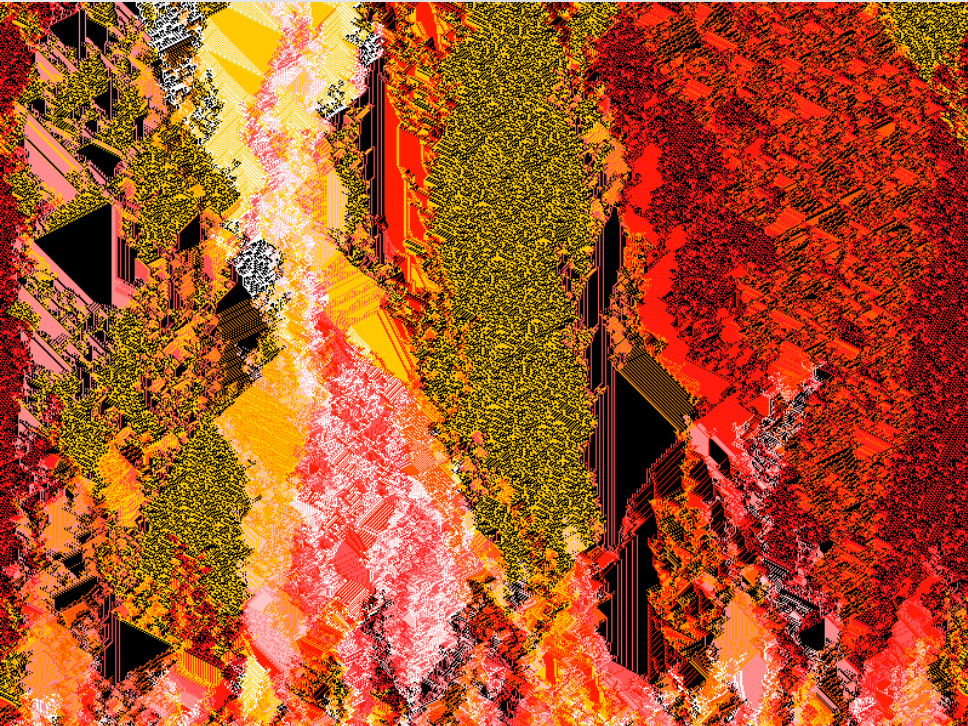
Examples



- $Q = \{0, 1\}$
- $V = \{-1, 0, 1\}$
- $f = \text{majority among neighbors}$

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CA as Dynamical Systems

► Cantor distance:

- $D(x, y)$ = dist. to center of the 1st cell where x and y differ
- Cantor distance: $d(x, y) = 2^{-D(x, y)}$
- the space of configurations is compact

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Theorem (Curtis, Hedlund, Lyndon, 69)

CA global functions are exactly the continuous function which commute with shift maps.

- shift map of vector z : $\sigma_z(x) = z' \mapsto x_{z+z'}$
- **intuition:** continuity \equiv short-term predictability

Deterministic Chaos

- δ : precision on initial conditions
- ϵ : desired long term precision



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Sensitivity to initial conditions:

$$\exists \epsilon, \forall x, \forall \delta, \exists y, \exists n : d(x, y) \leq \delta \text{ and } d(F^n(x), F^n(y)) \geq \epsilon$$



Deterministic Chaos

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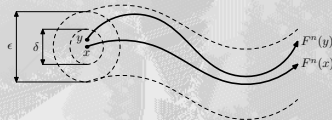


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Equicontinuity point at x :

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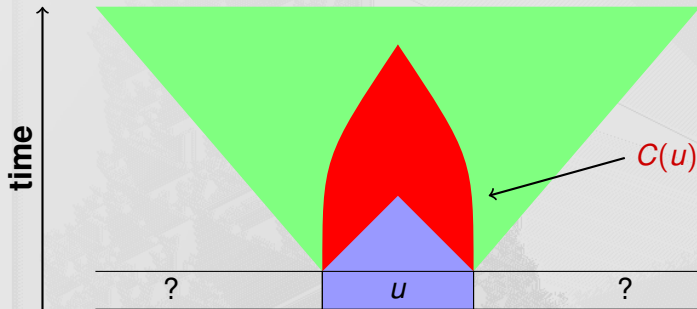
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- **Eq** $\stackrel{def}{=}$ set of CA having equicontinuity points

Equicontinuity without topology

Consequences $C(u)$ of a word u



$$[u] = \{x \in Q^{\mathbb{Z}} : x_0 \cdots x_{|u|-1} = u\}$$

$$C(u) = \{(z, t) : \forall x, y \in [u], F^t(x)_z = F^t(y)_z\}$$

Equicontinuity without topology

Walls

- F a 1D CA with radius r
- Obstacle $\stackrel{\text{def}}{=}$ word with infinite vertical strip in its consequences



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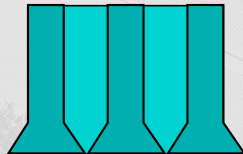
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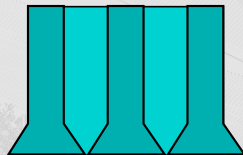
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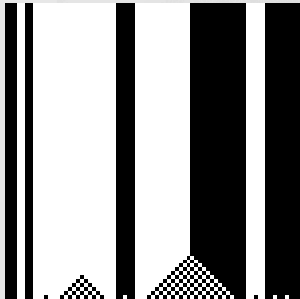
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Theorem (P. Kůrka, 1997)

$F \in \mathbf{Eq} \iff F \text{ has a wall} \iff F \text{ is almost everywhere equicontinuous}$

Examples: do they have walls?

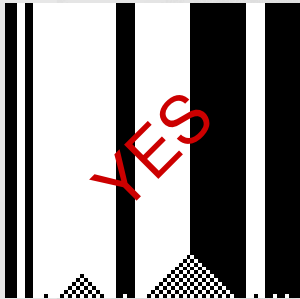


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Undecidability

The Usual Refrain

Theorem (Durand-Formenti-Varouchas, 2003)

It is undecidable to know whether a given CA has a wall or not.

(I've been told a short proof of this using Kari's undecidability result on nilpotency)

Corollary

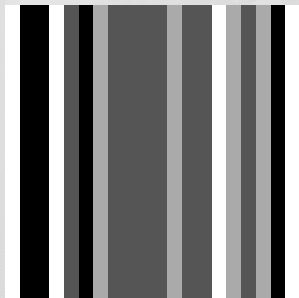
There is no computable bound on the size of the smallest wall.

Open problem

Is the set of 1D CA having a wall a Σ_2 -complete set?

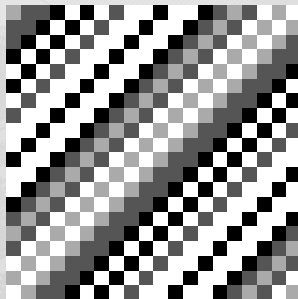
Problem with Shift Invariance

Identity



equicontinuous
long-term predictable

Shift



sensitive
NOT long-term predictable?

Topological dynamics with Cantor metrics not well adapted!

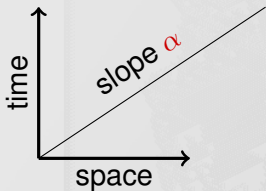
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The Work of M. Sablik (2008)



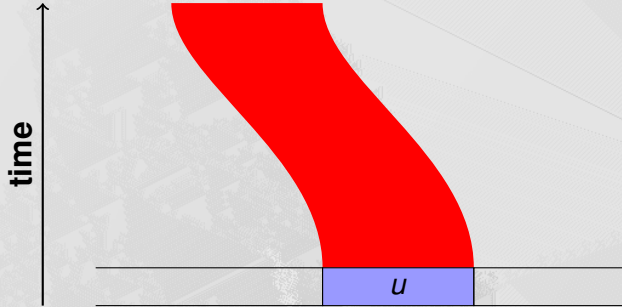
- 1 for any α , consider the 'sloped' system: $t \mapsto \sigma_{\lceil \alpha t \rceil} \circ F^t$
- 2 for each fixed α , Kůrka's results hold
- 3 study the set of slopes α showing a given behavior

Example of Theorem

For any given CA, the set of slopes with equicontinuity points is an interval of reals.

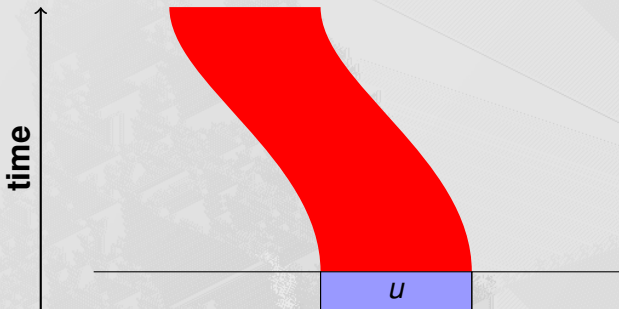
Generalization

- Walls along arbitrary curves



Generalization

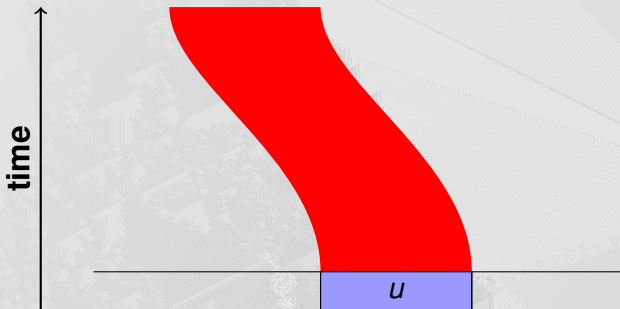
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- theory still works (wall \leftrightarrow equicontinuity) but...
- useful definition? do 'non-linear walls' exist?

Generalization

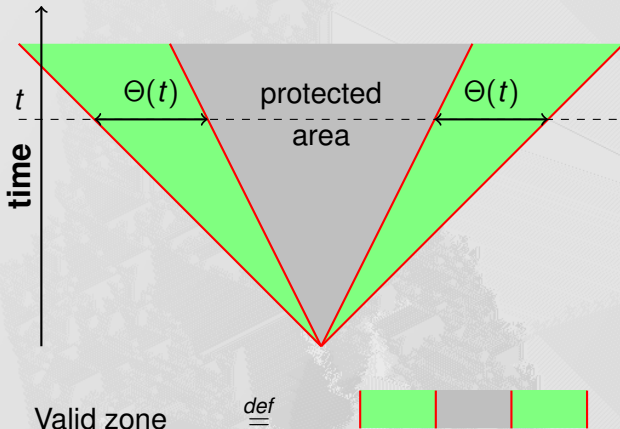
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- theory still works (wall \leftrightarrow equicontinuity) but...
- **useful definition? do 'non-linear walls' exist?**
- this talk: **YES** + other strange behaviors

The Key Construction

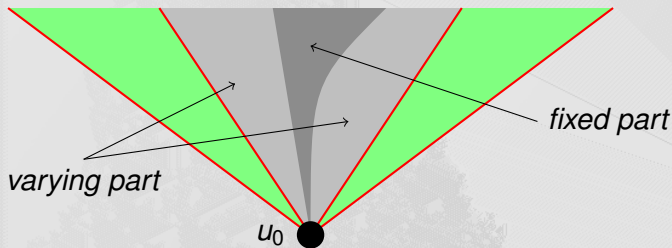
Outline



- 1 only a valid zone can erase a valid zone
- 2 after t steps, the age of any valid zone is at least t
- 3 when two valid zones meet, the older is destroyed

The Key Construction

Evil Twins



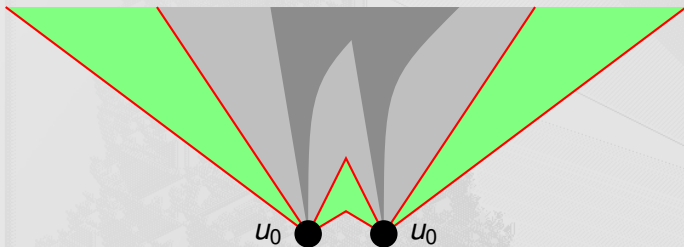
■ $u_0 \stackrel{\text{def}}{=} \text{protected zone of age 0}$

Principle

If the construction can merge with itself **then** $C(u_0)$ is exactly the fixed part of the construction.

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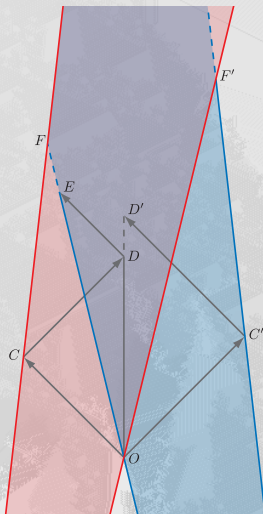
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Details



Result 1

Parabolic Consequences

Theorem

There exists a CA having a wall along a parabolic curve, but without any wall along any linear direction.

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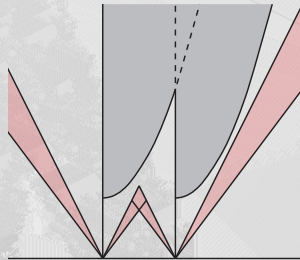
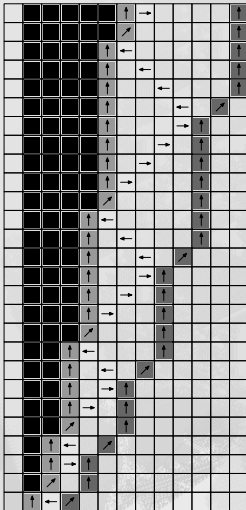
Main ideas

- two constructions
- intersection via Cartesian product
- protected zones



Result 1

2 pictures about the proof



Result 2

Characterization of linear slopes

Definition

A real α is computably enumerable (c.e.) if there is a computable sequence of rationals converging to it.

- left c.e. if the sequence is increasing
- right c.e. if the sequence is decreasing

Result 2

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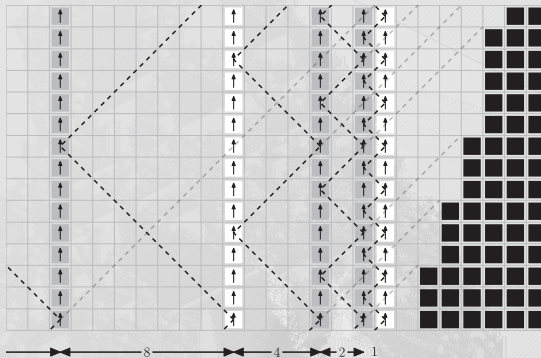
- **left** c.e. if the sequence is **increasing**
- **right** c.e. if the sequence is **decreasing**

Theorem

$[\alpha, \beta]$ is the set of slope of equicontinuity for some CA
iff α is **lce** and β is **rce**.

Result 2

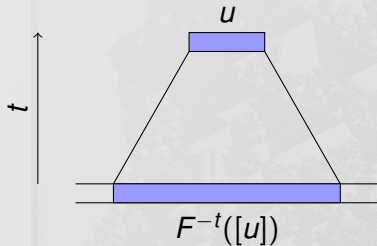
1 picture about the proof



$$\alpha = \sum_i \frac{\alpha_i}{2^i} \quad (\text{here: } \alpha_i = 0 \text{ iff } i = 0 \text{ or } i = 2)$$

Result 3

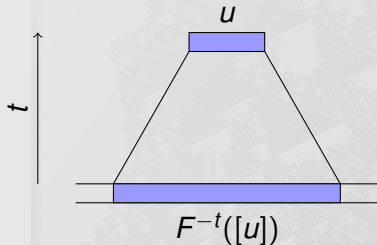
Strange Typical Behaviors



- μ : Bernoulli measure
- $\mu_t(u) \stackrel{\text{def}}{=} \mu(F^{-t}([u]))$
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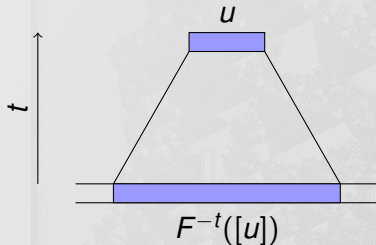
Proposition

There exist a CA F with:

- 1 $\mu_t(0) \rightarrow_t \chi$
- 2 $\mu_t(1) \rightarrow_t 1 - \chi$
- 3 $\mu_t(q) \rightarrow_t 0$ for any $q \notin \{0, 1\}$

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Strange Typical Behaviors



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■ u persistent $\stackrel{\text{def}}{\iff} \mu_t(u) \not\rightarrow 0$

Work in progress (with Sablik *et al.*)

Any transitive regular language can be the persistent language of some CA.

Other results + many details



Directional Dynamics along Arbitrary Curves in Cellular Automata

Delacourt, Poupet, Sablik, Theyssier

37p. To appear in TCS.

Future Work

- directional expansivity
- restriction to surjective CA
- what sets of configurations can be a μ -limit-set?
- what sets of configurations cannot be a μ -limit-set?