

Clandestine Simulations in Cellular Automata

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P. Guillon, P.-E. Meunier and *G. Theyssier*

CNRS/University of Turku,
Universidad de Chile/Université de Savoie,
CNRS/Université de Savoie

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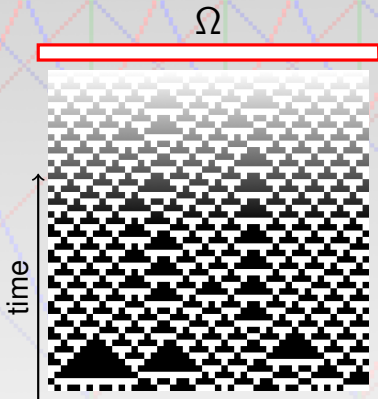
- 1** Limit sets and column factors
- 2** Sub-automata and factors
- 3** A word about proofs
- 4** Open problems

1 Limit sets and column factors

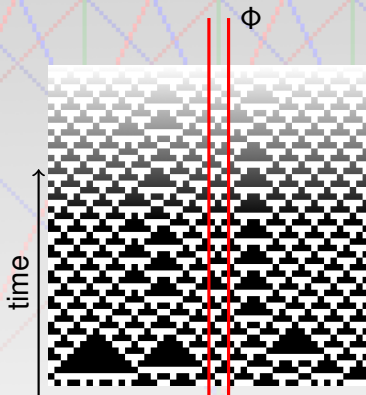
2 Sub-automata and factors

3 A word about proofs

4 Open problems



Limit set
Asymptotic observation



Column factor
Unprecise observation

- F a cellular automaton

- F a cellular automaton

- **limit set:** $\Omega_F = \bigcap_t F^t(Q^{\mathbb{Z}})$

- *upper-bound:* co-r.e. language

- *varying complexity:* from **finite** to **non-recursive**

- F a cellular automaton

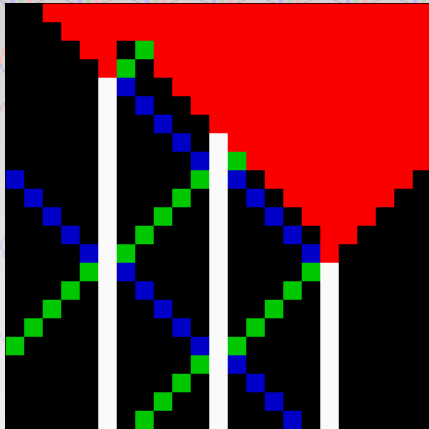
- **limit set:** $\Omega_F = \bigcap_t F^t(Q^{\mathbb{Z}})$

- *upper-bound:* co-r.e. language
- *varying complexity:* from **finite** to **non-recursive**

- **column factor of width k :** $\Phi_F = ([F^t(Q^{\mathbb{Z}})]_k)_t$

- *upper-bound:* context-sensitive language
- *varying complexity:* from **finite** to **non-context-free**

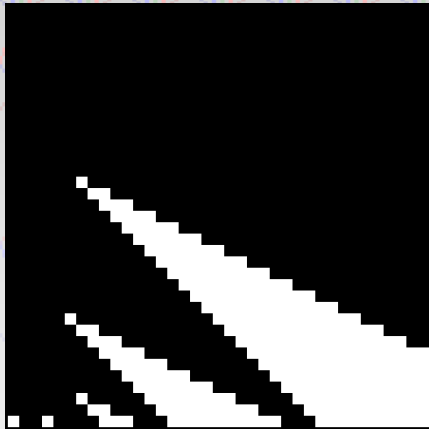
Example I



- left/right particles
- **symmetric** bounce on walls
- spreading state otherwise

Exercise

Show that it has a non-sofic limit set.



$$F(x)_i = x_{i+1}x_{i+2}$$

(example from P. Kůrka)

Exercise

Show that it has a non-sofic column factor.

Overview of the talk

1 Limit sets and column factors

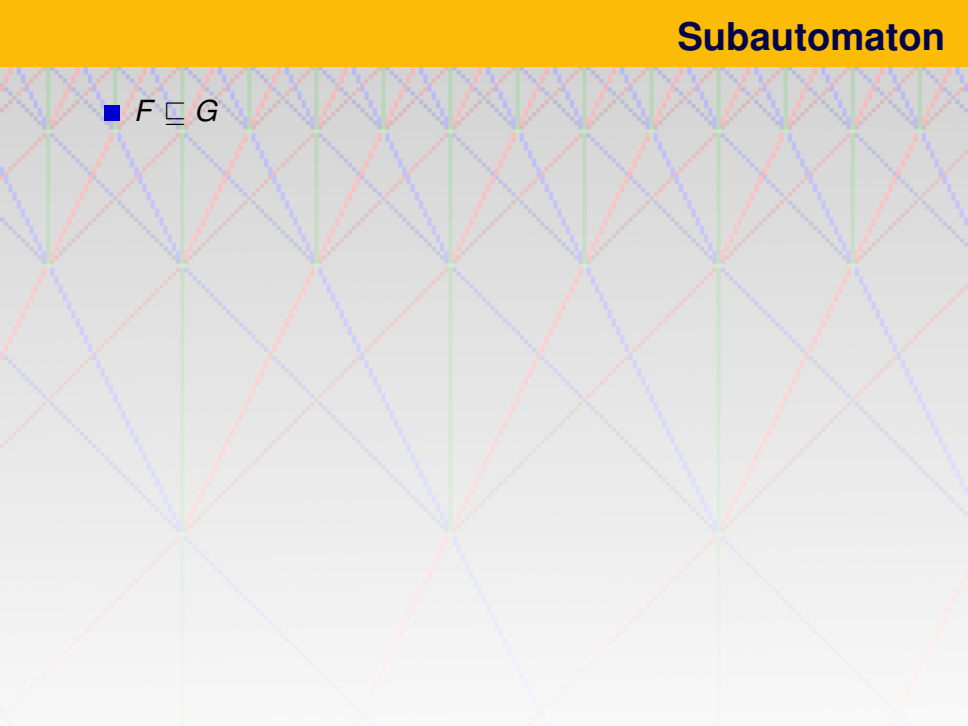
2 Sub-automata and factors

3 A word about proofs

4 Open problems

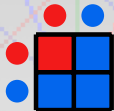
Subautomaton

■ $F \sqsubseteq G$

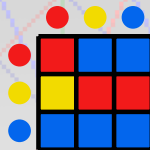


Subautomaton

■ $F \sqsubseteq G$



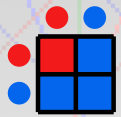
F



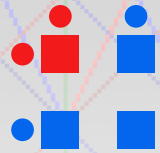
G

Subautomaton

■ $F \sqsubseteq G$



F

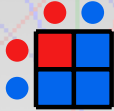


G

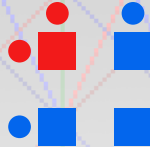
injection $\iota : Q_F \rightarrow Q_G$ with $\iota \circ F = G \circ \iota$

Subautomaton

■ $F \sqsubseteq G$



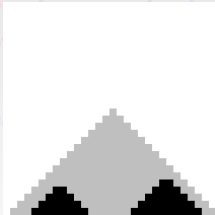
F



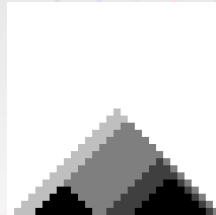
G

injection $\iota : Q_F \rightarrow Q_G$ with $\iota \circ F = G \circ \iota$

■ example:



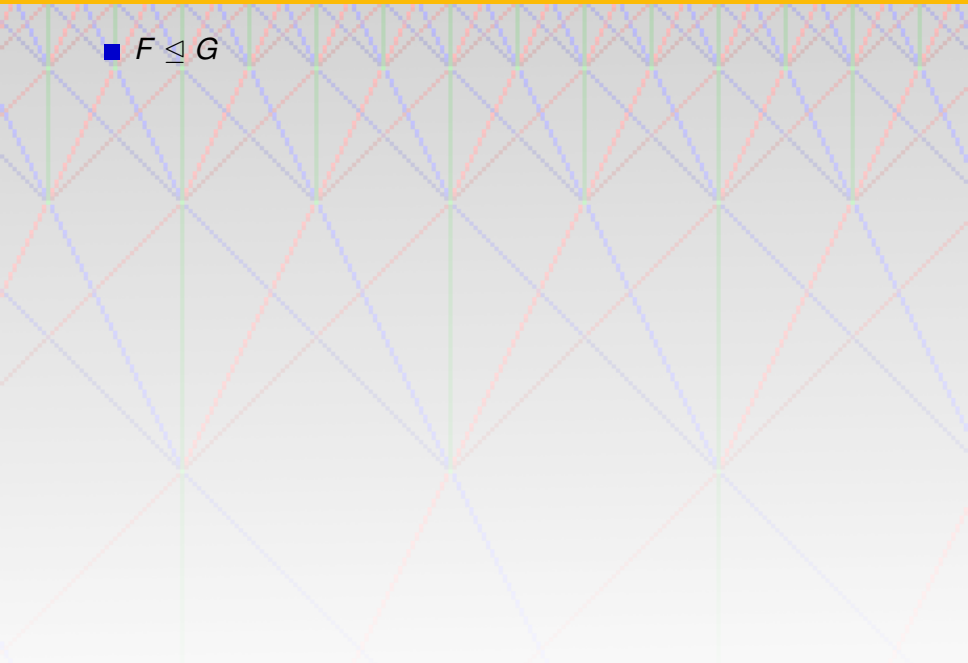
MAX with 3 states



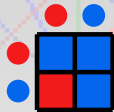
MAX with 5 states

Local Factor

■ $F \trianglelefteq G$



■ $F \trianglelefteq G$

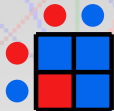


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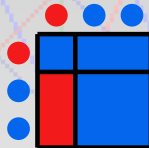


G

■ $F \trianglelefteq G$



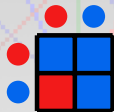
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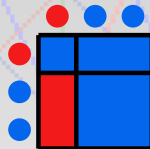
G

surjection $\pi : Q_G \rightarrow Q_F$ with $\pi \circ G = F \circ \pi$

■ $F \trianglelefteq G$



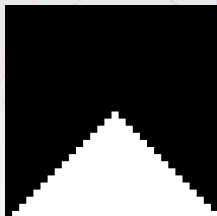
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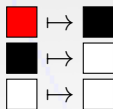
G

surjection $\pi : Q_G \rightarrow Q_F$ with $\pi \circ G = F \circ \pi$

■ example



MAX with 2 states



54 + spreading state

■ If $F \trianglelefteq_{\pi} G$ then (compactness argument)

$$\Omega_F = \pi(\Omega_G)$$

$$\Phi_F = \pi(\Phi_G)$$

► ***complexity increases from F to G***

- If $F \trianglelefteq_{\pi} G$ then (compactness argument)

$$\Omega_F = \pi(\Omega_G)$$

$$\Phi_F = \pi(\Phi_G)$$

- ▶ *complexity increases from F to G*

- If $F \sqsubseteq G$ then

$$\Omega_F \subseteq \Omega_G$$

$$\Phi_F \subseteq \Phi_G$$

- ▶ *does complexity increase from F to G ?*

Clandestine simulation

Complexity can be hidden in subautomata

The background of the slide features a complex, repeating geometric pattern. It consists of a grid of light gray squares. Overlaid on this grid are several sets of parallel lines in different colors: green, blue, and red. The green lines are vertical. The blue lines are diagonal, sloping downwards from left to right. The red lines are also diagonal, sloping upwards from left to right. The intersection of these lines creates a dense, woven appearance, with some lines being solid and others appearing as dashed or fainter lines, suggesting a layered or multi-dimensional structure. This visual metaphor likely represents the 'complexity hidden in subautomata' mentioned in the text.

Complexity can be hidden in subautomata

Theorem

Let F be any CA, then

- 1 there exists G with $F \sqsubseteq G$ and Φ_G **of finite type**
- 2 there exists G with $F \sqsubseteq G$ and Ω_G **NLOGSPACE-rec.**

Clandestine simulation

Complexity can be hidden in subautomata

Theorem

Let F be any CA, then

- 1 there exists G with $F \sqsubseteq G$ and Φ_G **of finite type**
- 2 there exists G with $F \sqsubseteq G$ and Ω_G **NLOGSPACE-rec.**

Corollary

There are intrinsically universal CA with a 'simple' limitset or a 'simple' column factor.

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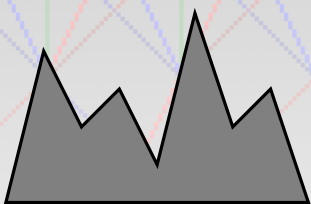
General idea: flooding



Resia Lake, Italy

General idea: flooding

■ $F \sqsubseteq G$

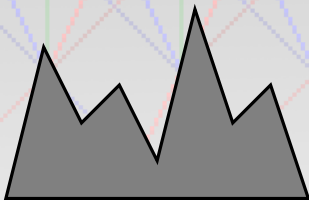


F

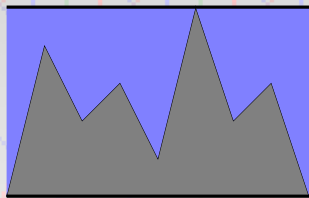
observation \equiv surface

General idea: flooding

■ $F \sqsubseteq G$



F



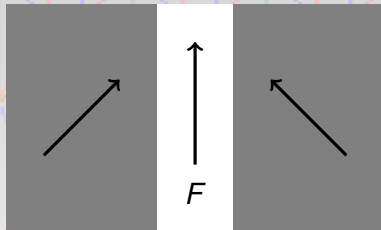
G

observation \equiv *surface*

complexity(F) \geq complexity(G)

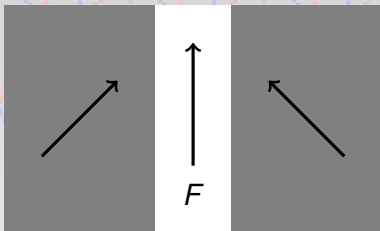
The case of column factor

- $G \approx F \times \text{transport layer}$



The case of column factor

- $G \approx F \times \text{transport layer}$



2-approximation of Φ_G

Set of columns which are valid when looking through a window of 2 time steps.

- **Fact:** Φ_G is equal to its 2-approximation
- **Corollary:** Φ_G is an SFT of order 2

The case of limit set

■ $G \approx F \cup (F \times S)$

- 1 on $F \times S$ component:** identity on F and **firing squad** on S
- 2 when S fires:** switch to F component
- 3 on F component:** do like F

The case of limit set

- $G \approx F \cup (F \times S)$

- 1 **on $F \times S$ component:** identity on F and **firing squad** on S

- 2 **when S fires:** switch to F component

- 3 **on F component:** do like F

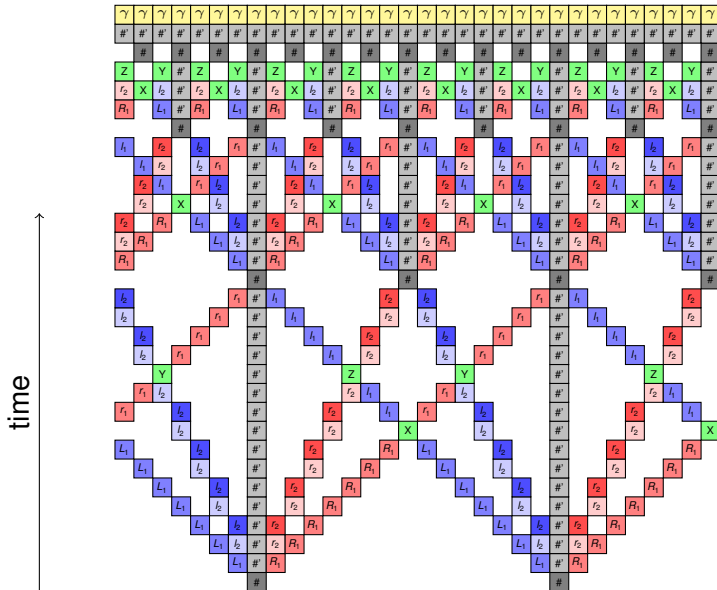
- **Fact:** the language of Ω_G is NLOGSPACE

- **easy part:** the whole F component is in Ω_G

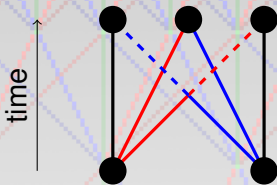
$$\Omega_G \cap Q_F^{\mathbb{Z}} = Q_F^{\mathbb{Z}}$$

- **hard part:** what else is in Ω_G ?

Jarkko Kari's Firing squad

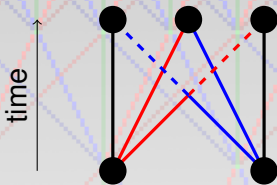


Euclidean reasoning



- valid $\stackrel{def}{=} \text{no fire and no spreading state}$
- valid discrete diagram $D \Rightarrow$ valid Euclidean diagram E with

$$E|_{\mathbb{Z}^2} = D$$



- valid $\stackrel{def}{=} \text{no fire and no spreading state}$
- valid discrete diagram $D \Rightarrow$ valid Euclidean diagram E with

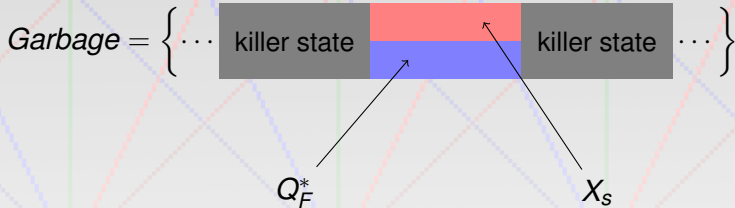
$$E|_{\mathbb{Z}^2} = D$$

■ Proof technique:

- 1 $u \in \Omega_G?$
- 2 try to construct an Euclidean history of u
- 3 get an Euclidean contradiction
- 4 conclude that $u \notin \Omega_G$

Garbage collection

$$\Omega_G = Q_F^Z \cup \textit{Garbage}$$



- X_S : confs with infinite history without killer or fire
- **(technical!) fact:** X_S is NLOGSPACE recognizable

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- Intrinsic universal CA:
 - with a SFT limit set?
 - with an injective limit system?
- is there a firing squad CA with an SFT limit set?
- factor maps:
 - is injectivity preserved by factor maps?
 - is there an expansive CA factor of a non-expansive CA?