

Amalgamation of Cellular Automata

JAC 2008 — *Médoc* session

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Motivations

- ▶ **The real one:** how common is (intrinsic) universality in CA?
- ▶ **Others:**
 - 1** amalgamations as composition operations
 - state efficient algorithmic constructions
 - alternative to cartesian product
 - 2** sub-families of CA obtained by syntactical constructions
 - study “malleability” of universality
 - study density of various properties
 - a hope to gain on decidability

Setting of the talk

- dimension is 1 (not necessary for most results)
- **a neighbourhood is fixed** (Von Neumann radius r)
- state set may vary but is always of the form $\{1, \dots, n\}$

Preliminary definitions

► Stability

- given F over state set Q
- $Q_0 \subset Q$ is F -stable if $F(Q_0^{\mathbb{Z}}) \subseteq Q_0^{\mathbb{Z}}$
- the induced automaton is denoted by $F|_{Q_0}$

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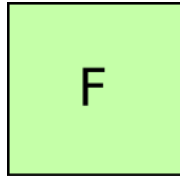
► Sub-automaton

- given F and G over Q_F and Q_G (resp.)
- **notation:** $F \sim G$ if they are equal upto state renaming
- **definition:** $F \sqsubseteq G$ (sub-automaton) if

$$\exists Q \subseteq Q_G : Q \text{ is } G\text{-stable and } G|_Q \sim F$$

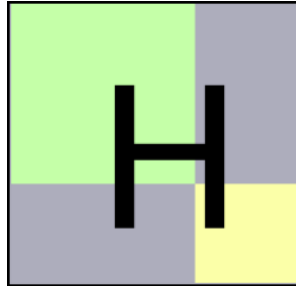
Amalgamations

- consider two CA:
 - F over Q_F
 - G over Q_G
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- H is an amalgamation of F and G if:
 - $Q_H = Q_F \cup Q_G$
 - Q_F and Q_G are both H -stable
 - $H|_{Q_F} = F$
 - $H|_{Q_G} = G$

Amalgamations

Operations & families

► **Amalgamation operation:**

■ $\Gamma : CA \times CA \rightarrow 2^{CA}$

■ $\Gamma(F, G)$ is a set of amalgamations of F and G

Amalgamations

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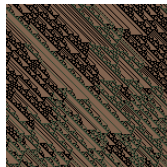
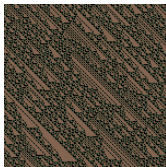
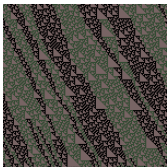
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► **Amalgamation family:**

- \mathbb{G} a finite set of CA
- Γ an amalgamation operation
- associated family $\mathcal{F}(\Gamma, \mathbb{G})$ is the closure of \mathbb{G} by Γ

Amalgamations

An example: scarf CA



Amalgamations

An example: scarf CA

- given F_1 and F_2
- choose:
 - 1 G with 2 states
 - 2 $\phi_1 : Q_{F_1} \cup Q_{F_2} \rightarrow Q_{F_1}$ (identity over Q_{F_1})
 - 3 $\phi_2 : Q_{F_1} \cup Q_{F_2} \rightarrow Q_{F_2}$ (identity over Q_{F_2})
- the **scarf** amalgamation H is defined by

$$H(x_{-r}, \dots, x_r) = \begin{cases} F_1(\phi_1(x_{-r}), \dots, \phi_1(x_r)) & \text{if } G(\chi(x_{-r}), \dots, \chi(x_r)) = 1 \\ F_2(\phi_2(x_{-r}), \dots, \phi_2(x_r)) & \text{if } G(\chi(x_{-r}), \dots, \chi(x_r)) = 2 \end{cases}$$

where $\chi(x) = i \iff x \in Q_{F_i}$.

Asymptotic density

- \mathcal{F} is a family (a set of CA with above conditions)
- \mathcal{P} is a property (another set)

► Density of \mathcal{P} in \mathcal{F} :

$$d_{\mathcal{F}}(\mathcal{P}) = \lim_{n \rightarrow \infty} \frac{\#\mathcal{F}_n \cap \mathcal{P}}{\#\mathcal{F}_n}$$

- \mathcal{F}_n is the set of CA from \mathcal{F} with n states
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Fact

If \mathcal{F} is an amalgamation family, then $d_{CA}(\mathcal{F}) = 0$

Simulation and universality

► Simulation:

- **intuition:** sub-automaton relation up to space-time rescaling
- **definition:** $F \preceq G$ iff

$$F \preceq G \iff \exists t, t', m, m' : \mathbf{b}_m \circ F^t \circ \mathbf{b}_m^{-1} \sqsubseteq \mathbf{b}_{m'} \circ G^{t'} \circ \mathbf{b}_{m'}^{-1}$$

where $\mathbf{b}_m : Q^{\mathbb{Z}} \rightarrow (Q^m)^{\mathbb{Z}}$ is the canonical bijection.

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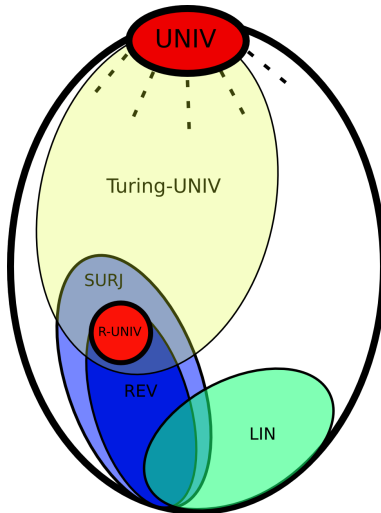
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► Intrinsic universality:

Definition

F is universal if for all G we have $G \preceq F$

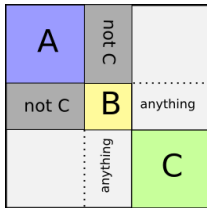
Simulation and universality



Amalgamations

Unambiguous families

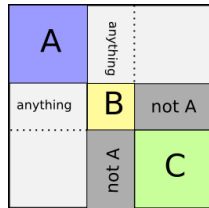
- an amalgamation operation is generally not associative



$\Gamma(\Gamma(A, B), C)$

\neq

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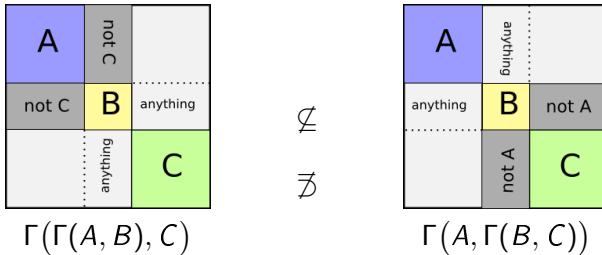


$\Gamma(A, \Gamma(B, C))$

Amalgamations

Unambiguous families

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- unambiguous amalgamation family $\mathcal{F}(\Gamma, \mathbb{G})$:

- 1 Γ is associative
- 2 $\#\Gamma(F, G)$ is a function of $|F|$ and $|G|$ only
- 3 there is no $G_1, G_2 \in \mathbb{G}$ with $G_1 \sqsubseteq G_2$

Amalgamations

A class of properties with high density

- $\mathcal{F}(\Gamma, \mathbb{G})$ an unambiguous amalgamation family
- a property \mathcal{P} is *malleable* if
 - 1 $\exists i, \forall G_1, \dots, G_i \in \mathbb{G} : \Gamma(G_1, \dots, G_i) \cap \mathcal{P} \neq \emptyset$
 - 2 \mathcal{P} is increasing for \sqsubseteq

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Proposition

If $\mathcal{F}(\Gamma, \mathbb{G})$ is unambiguous and \mathcal{P} is a malleable property then $d_{\mathcal{F}}(\mathcal{P}) = 1$.

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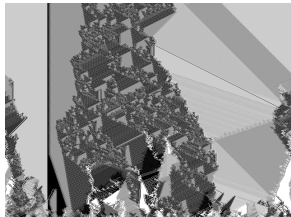
Proof ingredients:

- partition \mathcal{F}_n according to generator sequence
- balance of amalgamation application trees

Examples: Captive CA

► local constraint

$$(x_1, \dots, x_k) \mapsto y \in \{x_1, \dots, x_k\}$$

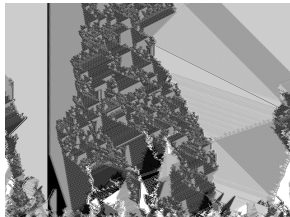


- $\Gamma_{\mathcal{K}}$: transition tables completed according to the constraint
- $\mathcal{K} = \mathcal{F}(\Gamma_{\mathcal{K}}, \{\mathbb{1}\})$

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Fact

- $\Gamma_{\mathcal{K}}$ is associative
- any \mathcal{P} increasing for \preceq and non-trivial in \mathcal{K} is malleable for $\mathcal{F}(\Gamma_{\mathcal{K}}, \mathbb{G})$

Examples: Majority CA

► local constraint

$$(x_1, \dots, x_k) \mapsto y \in \{x_i : c(i) \geq c(j), \forall j\}$$

where $c(i) = \#\{j : x_j = x_i\}$



- $\Gamma_{\mathcal{M}_{aj}}$: transition tables completed according to the constraint
- $\mathcal{M}_{aj} = \mathcal{F}(\Gamma_{\mathcal{M}_{aj}}, \{\mathbb{1}\})$

Fact

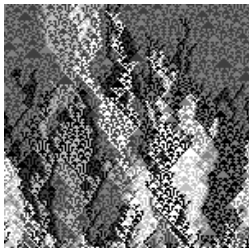
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- $\Gamma_{\mathcal{M}_{ino}}$: transition tables completed according to the constraint
- $\mathcal{M}_{ino} = \mathcal{F}(\Gamma_{\mathcal{M}_{ino}}, \{\mathbb{1}\})$

Fact

- $\Gamma_{\mathcal{M}_{ino}}$ is associative
- any \mathcal{P} increasing for \preceq and non-trivial in \mathcal{M}_{ino} is malleable for $\mathcal{F}(\Gamma_{\mathcal{M}_{ino}}, \mathbb{G})$

Encodings

Definition

- let \mathcal{F}_1 and \mathcal{F}_2 be families of CA
- $\phi : \mathcal{F}_1 \rightarrow \mathcal{F}_2$ is an encoding if
 - 1 it is computable and injective
 - 2 $\forall F \in \mathcal{F}_1 : F \preceq \phi(F)$
- it is *faithful* if $F \in \mathcal{U} \iff \phi(F) \in \mathcal{U}$

► Remarks

- an encoding preserves existence of universal CA
- a faithful encoding preserves undecidability of \mathcal{U}

Encodings

Results

Proposition

$$\begin{array}{ccccc} & & & \mathcal{M}_{aj} & \rightarrow & \mathcal{F}(\Gamma_{\mathcal{M}_{aj}}, \mathbb{G}) \\ & & \nearrow & & & \\ CA \Rightarrow \mathcal{K} & \rightarrow & \mathcal{M}_{ino} & \rightarrow & \mathcal{F}(\Gamma_{\mathcal{M}_{ino}}, \mathbb{G}) \\ & & \searrow & & & \\ & & & \mathcal{F}(\Gamma_{\mathcal{K}}, \mathbb{G}) & & \end{array}$$

\rightarrow : encoding

\Rightarrow : faithful encoding

► Remarks:

- $r \geq 5$ needed for “ $CA \Rightarrow \mathcal{K}$ ”
- $\mathbb{G} \notin \mathcal{U}$ needed for “ $\mathcal{K} \Rightarrow \mathcal{F}(\Gamma_{\mathcal{K}}, \mathbb{G})$ ”

Encodings

Corollaries

Assembling previous results we have:

- ▶ $d_{\mathcal{F}}(\mathcal{U}) = 1$ for the following \mathcal{F} :
 - \mathcal{K} , any unambiguous $\mathcal{F}(\Gamma_{\mathcal{K}}, \mathbb{G})$
 - \mathcal{M}_{aj} , any unambiguous $\mathcal{F}(\Gamma_{\mathcal{M}_{aj}}, \mathbb{G})$
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- ▶ for $\mathcal{F} = \mathcal{K}$ or $\mathcal{F} = \mathcal{F}(\Gamma_{\mathcal{K}}, \mathbb{G})$:
 - \mathcal{U} is undecidable
 - ...

Captive CA and reversibility

Proposition

For any RCA F there exists a RCA G with:

- $F \preceq G$
- G and G^{-1} have a common wall state w

$$\forall x_{-r}, \dots, x_r, G(x_{-r}, \dots, x_{-1}, w, x_1, \dots, x_r) = w$$

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► Idea: define G over $Q_F \times Q_F \cup \{w\}$ with



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Proposition

There exists a reversible-universal captive CA.

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There exists a reversible-universal captive CA.

► Proof sketch

- start from a reversible-universal CA F with a wall state
- use classical simulation of F over configuration of the form

$$\cdots \# q_1 \cdots q_n \# x_0 \# q_1 \cdots q_n \# x_1 \# q_1 \cdots q_n \# x_2$$

where $Q_F = \{q_1, \dots, q_n\}$

- local checking of this SFT condition
- interpret errors as the wall state

Unresolved problems

- faithful encoding for \mathcal{M}_{aj} and \mathcal{M}_{ino}
- $\Gamma_F(G, H) \cap \mathcal{U}$ when $\{F, G, H\} \cap \mathcal{U} = \emptyset$?
- a reversible universal captive CA in 2D?

Future directions

- extension to other amalgamation families
- decidability
- families for the reversible case:
 - with almost only reversible CA (2D case)?
 - with almost only reversible-universal CA?
- μ -limit sets
- the case of fixed state set and increasing neighbourhood

End of Médoc session...



Questions?

