

EMERGENCE OF BACKGROUND  
AND PARTICULE-LIKE STRUCTURES  
IN ONE-DIMENSIONAL CA

An Attempt of Formalization

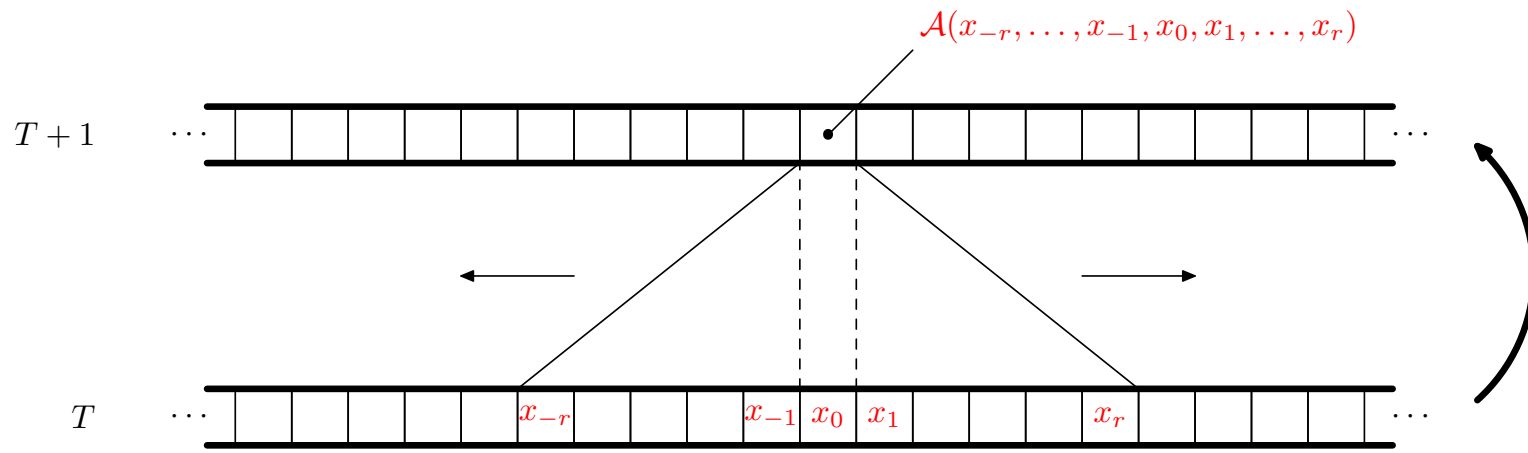
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*AUTOMATA 2002, 8<sup>th</sup> International Workshop on Cellular Automata,  
September 12-14, Prague.*

# INTRODUCTION AND BASIC OBSERVATIONS

# One-dimensional CA



Dual point of view

Parallel computing device

$$\mathcal{A} : A^{2r+1} \rightarrow A$$

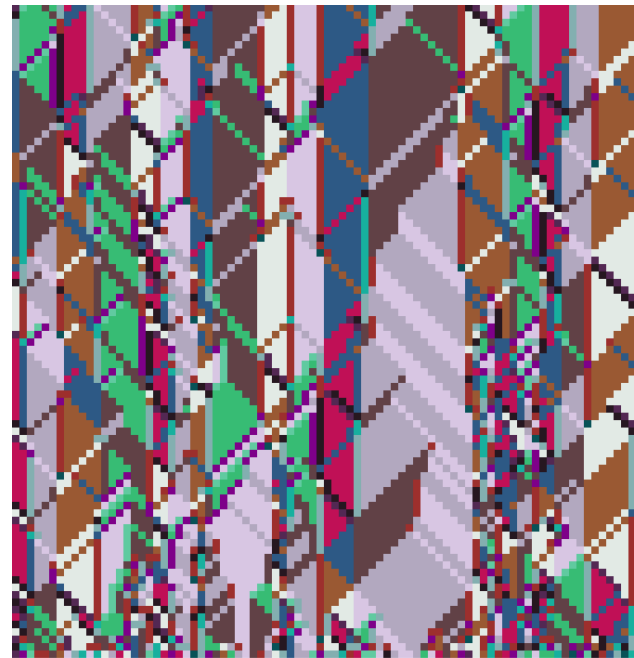
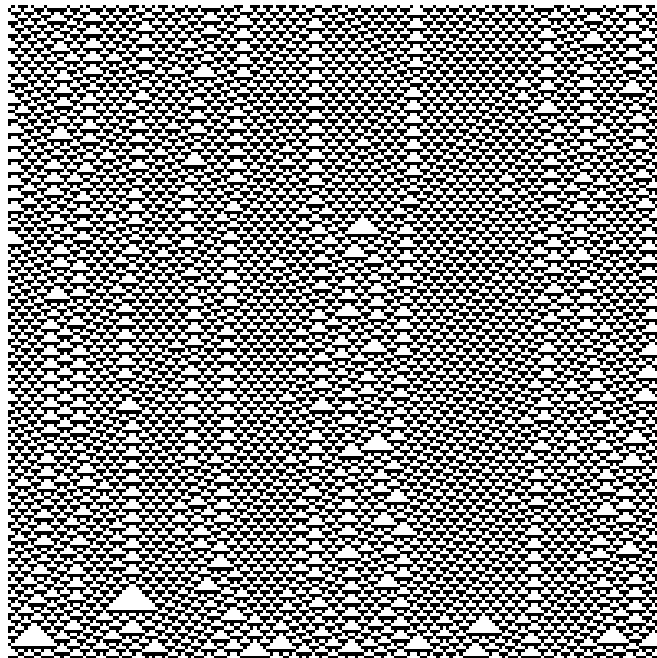
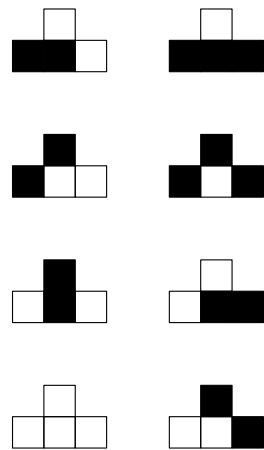
a **local** rule **uniformly** and  
**synchronously** applied to the  
line of cells

Discrete dynamical system

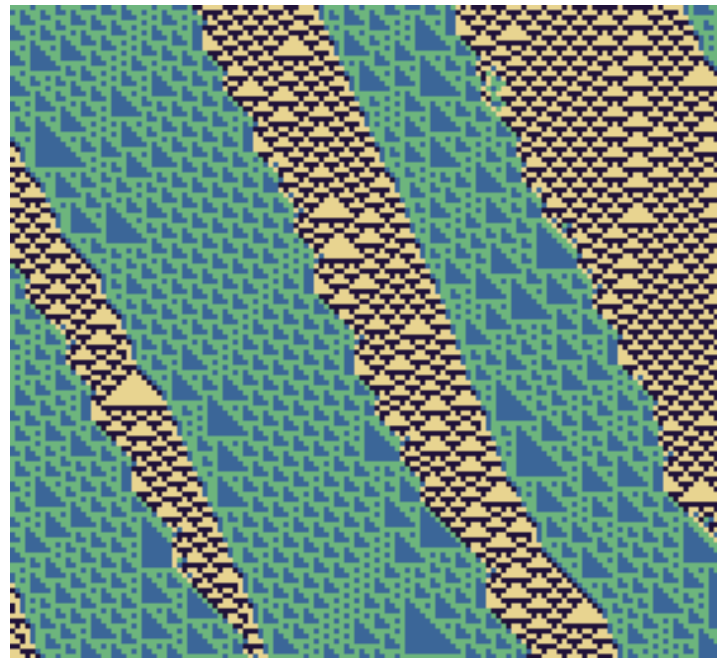
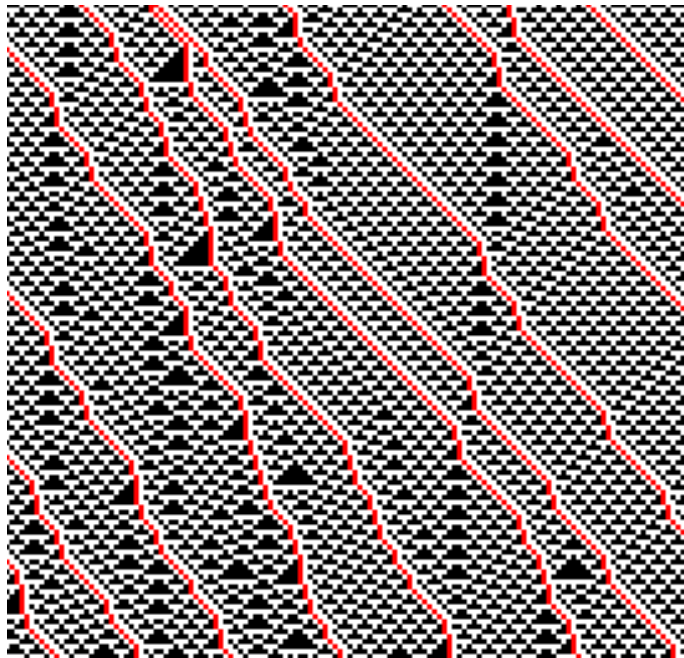
$$\mathcal{A} : A^{\mathbb{Z}} \rightarrow A^{\mathbb{Z}}$$

a **continuous** function which  
**commute with the shift** on the  
**compact** space  $A^{\mathbb{Z}}$

# Elementary CA $W54$



More complicated examples



Informal proposal for a generalized notion of background

Up to **re-scaling**, a **sub-automaton** which is ultimately **active** on a **sizeable part** of the line of cells.

# PRELIMINARY NOTIONS



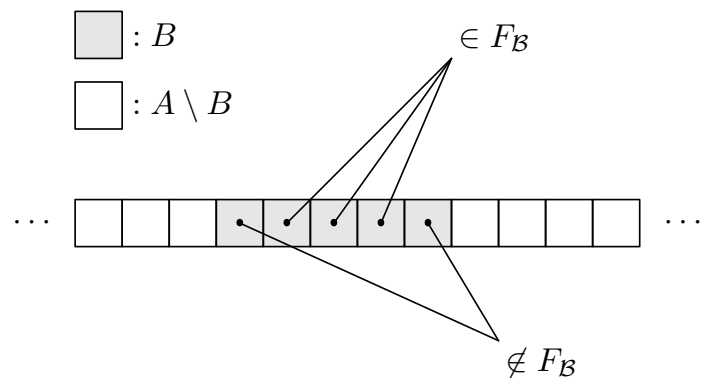
Sub-automata and “activity”

Definition :

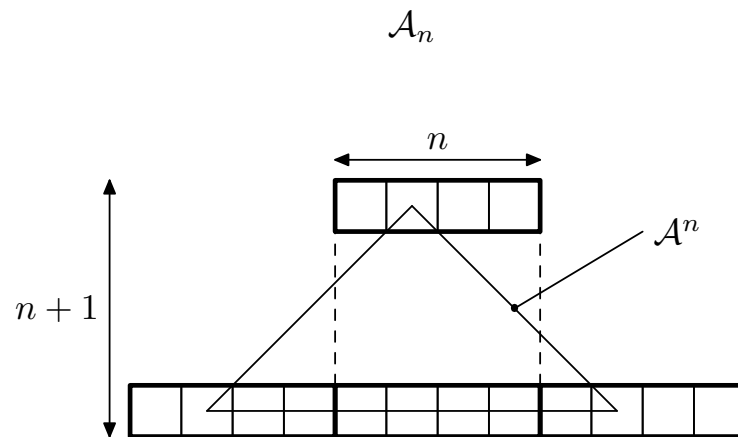
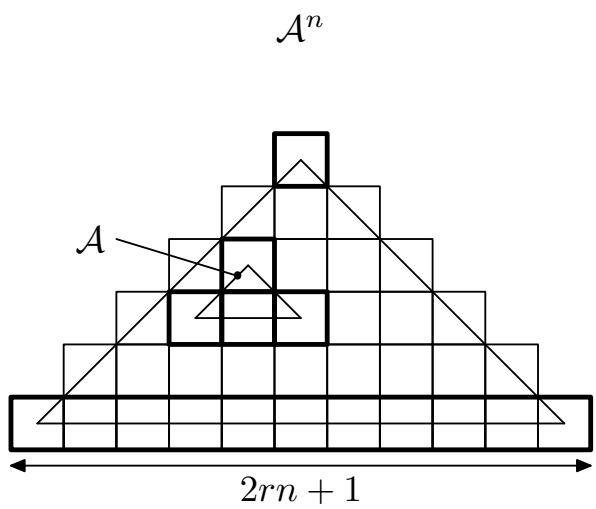
$\mathcal{B}$  is a **sub-automaton** of  $\mathcal{A}$  (written  $\mathcal{B} \subseteq \mathcal{A}$ ) if there is an injection  $\phi : B \rightarrow A$  s.t.  $\phi \circ \mathcal{B} = \mathcal{A} \circ \phi$  on  $B$

Definition :

Let  $\mathcal{B} \subseteq \mathcal{A}$  and  $c \in A^{\mathbb{Z}}$ . The set of **active sites** of  $\mathcal{B}$  in  $c$  is defined as  $F_{\mathcal{B}}(c) = \{i : \forall j, |j| \leq r, c_{i+j} \in B\}$



Re-scaling and grouping operation



Probabilistic tools

$A^{\mathbb{Z}}$  is endowed with the **uniform product measure** :  $\mu$ .

Definition :

The **density** of  $X \subseteq \mathbb{Z}$  is defined as

$$d(X) = \limsup_{n \rightarrow \infty} \frac{|X \cap \{-n, \dots, n\}|}{2n + 1}$$

# THE NOTION OF $\alpha$ -BACKGROUND

## Formal definition

## Definition :

For  $0 \leq \alpha \leq 1$ , a CA  $\mathcal{B}$  is an  **$\alpha$ -background** for a CA  $\mathcal{A}$  on  $\mathcal{D}$  (set of initial configurations) if  $\exists n \in \mathbb{N}_+$  s.t. :

- $\mathcal{B} \subsetneq \mathcal{A}_n$  ;
- $\forall c \in \mathcal{D}, \exists t_0 \in \mathbb{N}_+$  s.t.  $\forall t \geq t_0 : d(F_{\mathcal{B}}(\mathcal{A}_n^t(c))) \geq \alpha$ .

An  $\alpha$ -background is said to be **visible** if  $\mathcal{D}$  has a strictly positive measure.

About visibility

Definition :

Let  $\sigma$  be the **shift** automaton  $\sigma : c \mapsto c'$  with  $c'(i) = c(i - 1)$

Theorem :

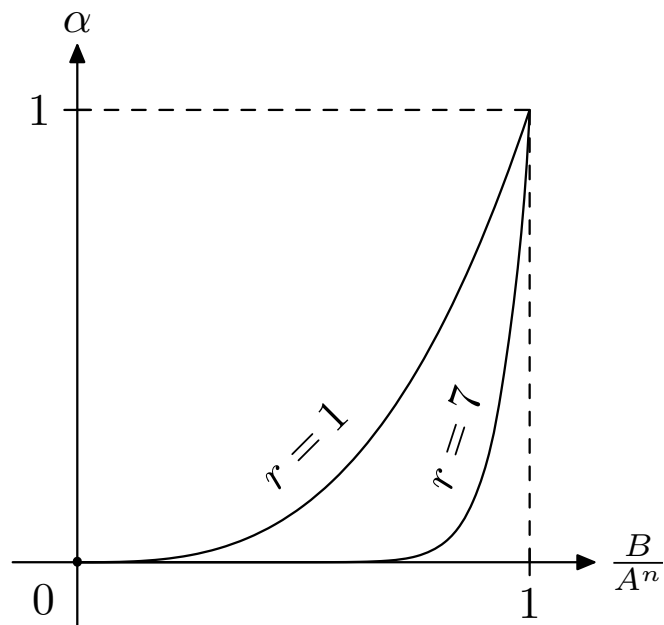
$\sigma$  is ergodic (*i.e.* any measurable set closed under  $\sigma$  is either of measure 1 or of measure 0).

Corollary :

**visible**  $\alpha$ -background  $\iff \mu(\mathcal{D}) = 1$ .

FIRST P PROPERTIES

## Density and surjectivity



Lemma :

If  $\mathcal{A}$  is **surjective** :

$\mathcal{B}$   $\alpha$ -background of  $\mathcal{A}$  of size

$$n \text{ implies } \alpha = \left( \frac{|B|}{|A|^n} \right)^{2r+1}$$



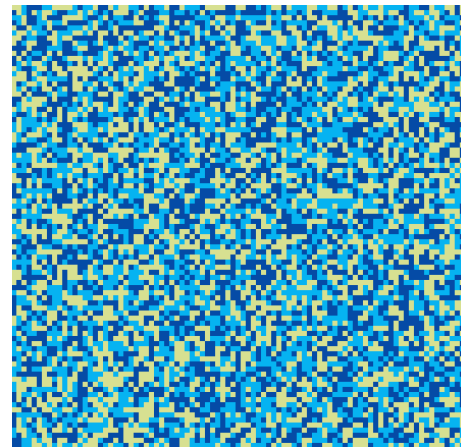
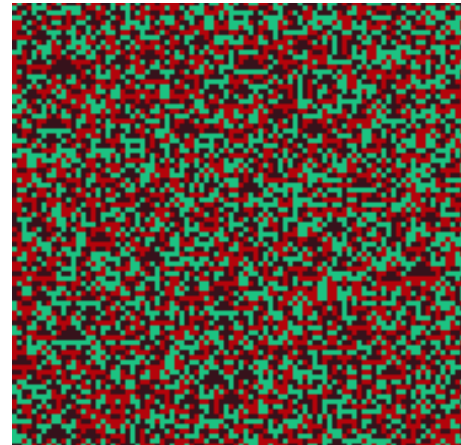
« Absence » of  $\alpha$ -background

Lemma :

An additive CA  $\mathcal{A}$  of radius  $r$  defined on  $\mathbb{Z}_p$  (with  $p$  prime  $p|2r + 1$ ) is such that  $\mathcal{B} \subseteq \mathcal{A}_n \Rightarrow \frac{|\mathcal{B}|}{|\mathcal{A}|^n} \leq \frac{1}{p}$ .

Corollary :

For arbitrarily small  $\alpha$  there is some CA with no visible  $\alpha$ -background.



Constructing arbitrary  $\alpha$ -backgrounds

## Theorem :

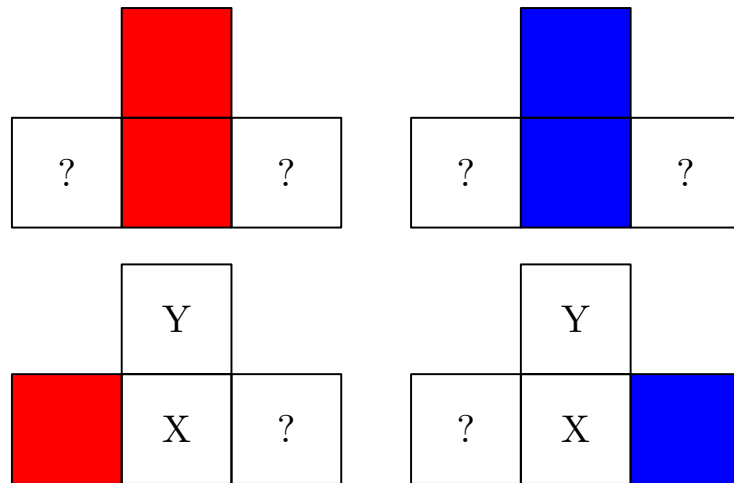
For any CA  $\mathcal{A}$ ,  $\exists$  CA  $\mathcal{A}'$  s.t.  $\mathcal{A} \subseteq \mathcal{A}'$  and  $\forall n : \mathcal{B} \subseteq \mathcal{A}_n \Rightarrow \mathcal{B}$  is a visible  $\alpha$ -background of  $\mathcal{A}'$  with  $\alpha > 0$ .

## Corollary :

$\exists$  CA with infinitely many visible  $\alpha$ -backgrounds (each with  $\alpha > 0$ ).

Proof sketch

$\forall n \in \mathbb{N}_+, \forall \mathcal{B} \subseteq \mathcal{A}_n$  build 2 wall states (left and right) such that :

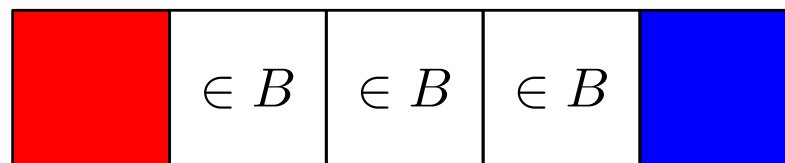


$$X \in B \Rightarrow Y \in B.$$

End of proof sketch

The wall states are obtained by adding a synchronizing component to  $\mathcal{A}$  (solution to the **firing squad** problem).

Then the following pattern ensures that  $\mathcal{B}$  is an  $\alpha$ -background for  $\mathcal{A}$  (the active sites for  $\mathcal{B}$  between the walls are preserved under iterations) :



Undecidability result

Definition :

A CA is **nilpotent** if it ultimately leads to the same stable configuration from any initial configuration.

Property :

A CA is nilpotent if and only if it admits the trivial 1-state CA as a visible 1-background.

Corollary :

It is undecidable to know if, given  $\mathcal{A}$  and  $\mathcal{B}$  and  $\alpha$ ,  $\mathcal{B}$  is a visible  $\alpha$ -background of  $\mathcal{A}$ .